

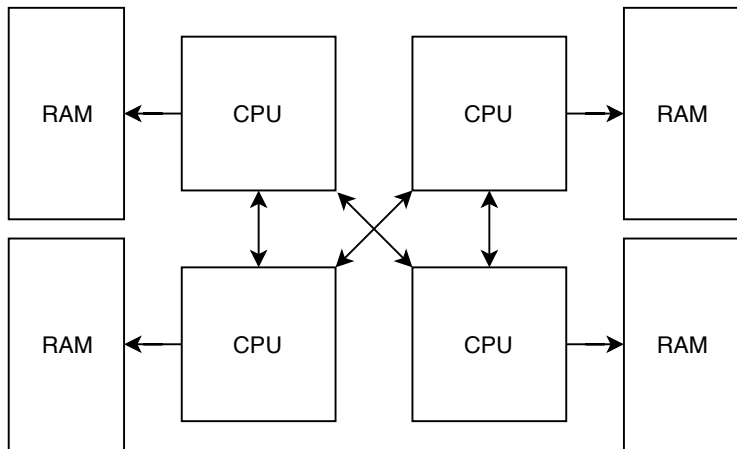
NUMA

Mark Johnston
markj@FreeBSD.org



FreeBSD Bay Area Vendor Summit
October 11, 2019

Non-Uniform Memory Access



OS Responsibilities

Minimize remote memory accesses

- ▶ Avoid remote access latency penalty
- ▶ Avoid bottlenecking on cross-domain interconnect

Requirements:

- ▶ Balance resource utilization
- ▶ Allow applications to provide hints (scheduling, memory allocation)
- ▶ Handle local memory shortages gracefully
- ▶ Affinitize static data structures



FreeBSD

APIs

Kernel:

- ▶ `bus_get_domain(9)`, `bus_dma_tag_set_domain(9)`
- ▶ `malloc_domainset(9)`, `kmem_malloc_domainset(9)`
- ▶ `uma_zalloc_domain(9)` (slow!)

Userspace:

- ▶ `cpuset(1)`
- ▶ `cpuset_getdomain(2)` `cpuset_setdomain(2)`



Review: Domain Selection Policies, domainset (9)

DOMAINSET_POLICY_ROUNDROBIN

- ▶ Cycle through domains: `d = iter++ % ds->ds_cnt`
- ▶ 0, 1, 2, 3, 0, 1, 2, 3, 0, ...

DOMAINSET_POLICY_FIRSTTOUCH

- ▶ Pick the domain of the current CPU: `d = PCPU_GET(domain)`
- ▶ Userland default, good for short-lived processes

DOMAINSET_POLICY_PREFER

- ▶ Pick the domain specified in the policy: `d = ds->ds_prefer`
- ▶ Fall back to round-robin when free pages are scarce

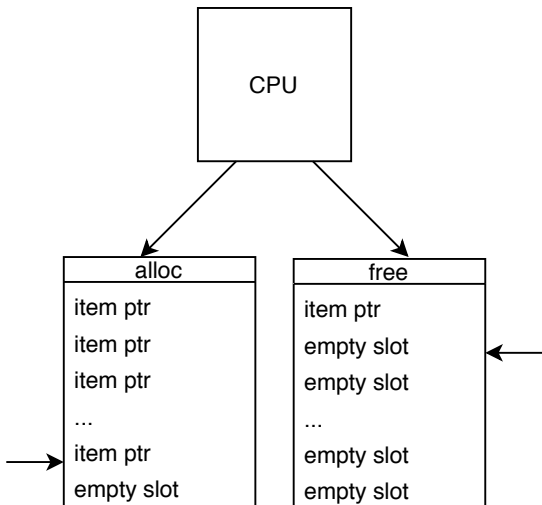
DOMAINSET_POLICY_INTERLEAVE

- ▶ Round-robin with a stride
- ▶ 0, 0, ..., 0, 1, 1, ..., 1, 0, 0, ...
- ▶ Superpage-friendly: use a stride of 512
- ▶ Kernel default



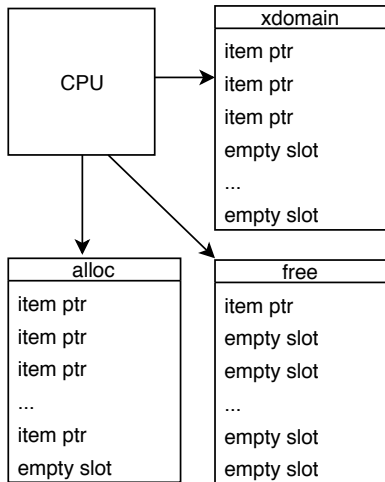
Review: UMA per-CPU caches

- ▶ Bucket: dynamically allocated array
- ▶ Items allocated from alloc bucket
- ▶ Items freed to free bucket
- ▶ Buckets are swapped if empty (alloc) or full (free)
- ▶ Per-domain cache of full buckets
- ▶ Slow path: lock the zone, check bucket cache

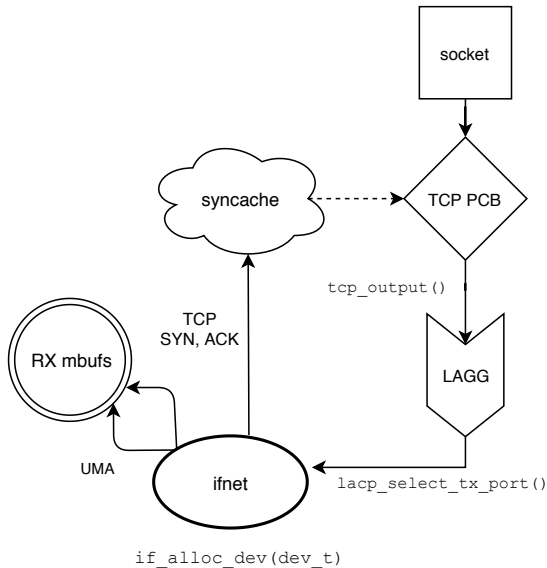


options UMA_XDOMAIN

- ▶ On free, find item domain
- ▶ Cache in free if domain == PCPU_GET(domain), else xdomain
- ▶ Slow path: lock the zone, drain xdomain
- ▶ Special optimization for 2 domains



Network affinity



vm_page_array (amd64 only)

- ▶ One vm_page structure per 4KB page
- ▶ vm_page_array allocated early during boot
- ▶ Physically contiguous → allocated from single domain
- ▶ Unfriendly to first-touch allocation policy
- ▶ Now backed by “correct” memory, up to 2MB boundaries



Other Data Structures

- ▶ PCPU area (amd64)
- ▶ ULE per-CPU thread queues
- ▶ callout wheel
- ▶ vm_page locks (by removing their usage)
- ▶ Kernel thread stacks



Memory-bound pgbench on a 2-socket system, r353116

Core	IPC	Instructions	Cycles	Local DRAM accesses	Remote DRAM Accesses
0	0.46	1097 M	2402 M	1467 K	759 K
1	0.45	1090 M	2402 M	1464 K	766 K
2	0.46	1095 M	2402 M	1556 K	801 K
3	0.46	1096 M	2402 M	1445 K	755 K
4	0.46	1099 M	2402 M	1507 K	787 K
5	0.45	1091 M	2402 M	1550 K	813 K
6	0.46	1099 M	2402 M	1482 K	785 K
7	0.45	1092 M	2402 M	1509 K	790 K
8	0.46	1100 M	2402 M	1469 K	771 K
9	0.45	1090 M	2402 M	1535 K	800 K
10	0.46	1094 M	2402 M	1585 K	830 K
11	0.45	1092 M	2402 M	1507 K	777 K
12	0.46	1099 M	2402 M	1481 K	776 K
13	0.46	1095 M	2402 M	1482 K	780 K
14	0.46	1094 M	2402 M	1535 K	793 K
15	0.45	1092 M	2402 M	1516 K	776 K
16	0.41	992 M	2402 M	796 K	1256 K
17	0.41	991 M	2402 M	763 K	1208 K
18	0.43	1040 M	2402 M	851 K	1365 K
19	0.43	1034 M	2402 M	860 K	1390 K
20	0.43	1042 M	2402 M	840 K	1332 K
21	0.43	1030 M	2402 M	852 K	1404 K
22	0.43	1035 M	2402 M	857 K	1392 K
23	0.43	1035 M	2402 M	836 K	1335 K
24	0.43	1039 M	2402 M	834 K	1341 K
25	0.43	1034 M	2402 M	830 K	1335 K
26	0.43	1040 M	2402 M	838 K	1339 K
27	0.43	1035 M	2402 M	841 K	1335 K
28	0.43	1040 M	2402 M	835 K	1321 K
29	0.43	1038 M	2402 M	818 K	1319 K
30	0.43	1041 M	2402 M	806 K	1269 K
31	0.43	1031 M	2402 M	831 K	1327 K



Future Direction

- ▶ Continue affinitizing static kernel data structures
 - ▶ e.g., `vm_reserv_array`, `vm_dom[]`
- ▶ Taskqueue affinity
- ▶ NUMA awareness in UMA by default
- ▶ Improve NUMA support on !amd64
- ▶ ...?