



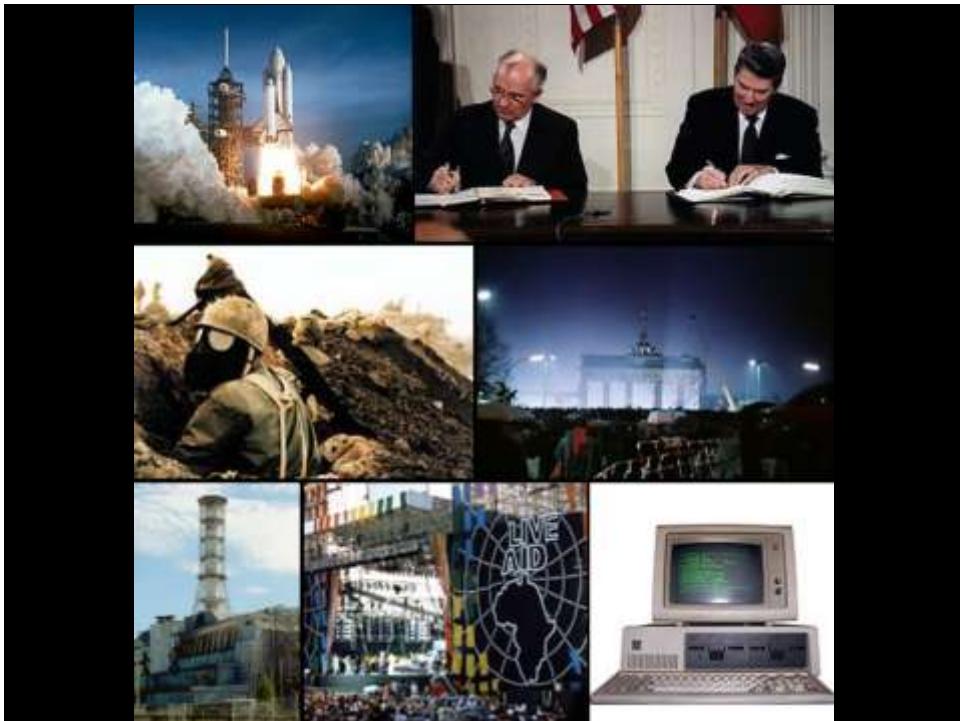
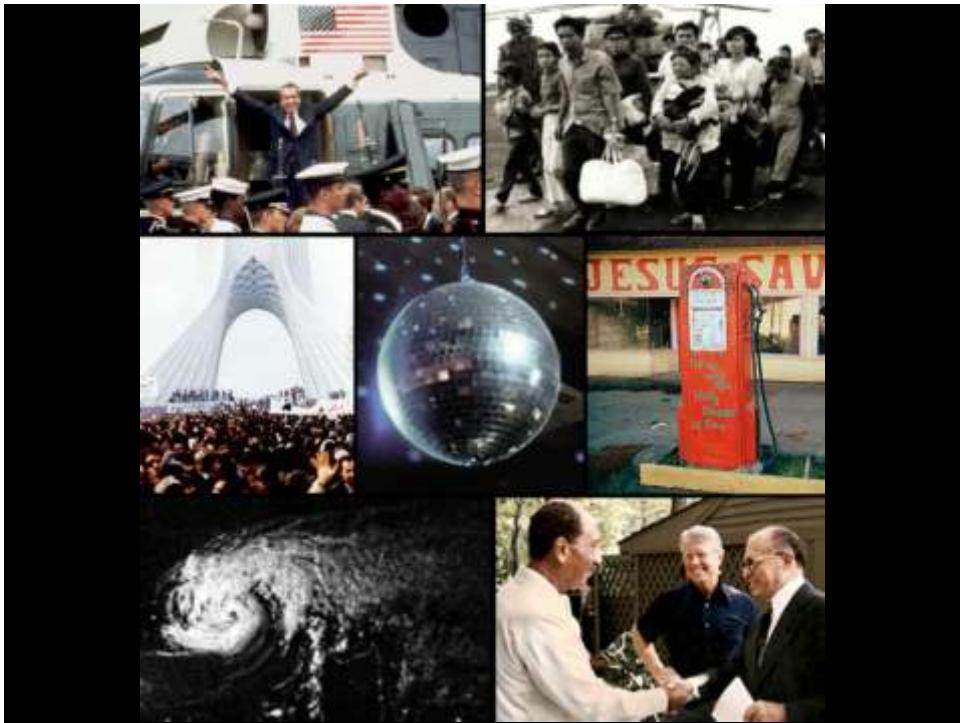
Unix Architecture Evolution: Milestones and Lessons Learned

Diomidis Spinellis

Department of Management Science and Technology
Athens University of Economics and Business

www.spinellis.gr
@CoolSWEng









Overview

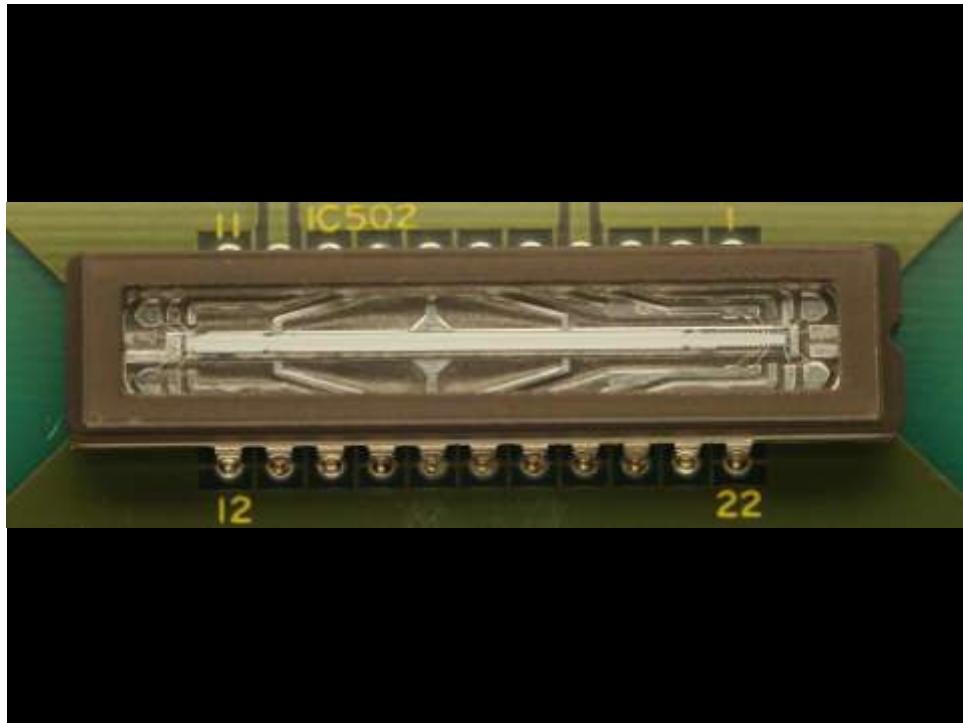
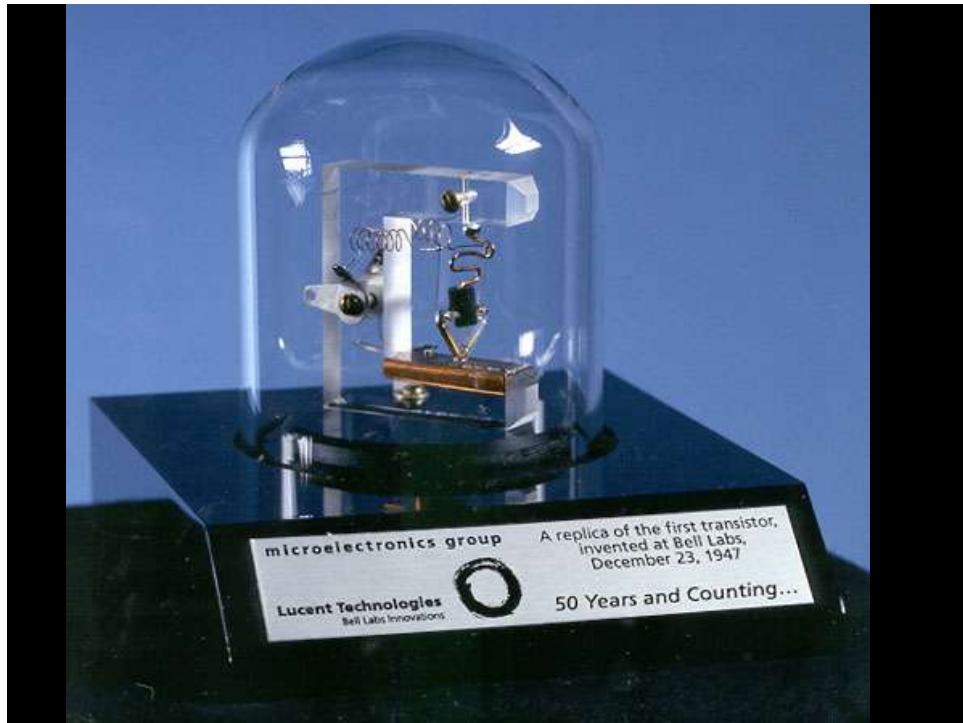
- Groundwork
- Sources
- Important architectural milestones
- Evolution in numbers
- Evolution in words











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Vol. 27, pp. 379-411, 1948, Bell Telephone Labs.

A Mathematical Theory of Communication

By C. E. SHANNON

INTRODUCTION

THREE decades of development of various methods of modulation such as PCM and PPM which exchange telephone signals-to-carry over long distances have introduced the interest in a general theory of communications. A basis for such a theory is contained in the important papers of Nyquist¹ and Hartley² on this subject. In the present paper we will extend the theory to include a number of new factors. In particular the effect of noise on the transmission and reception of messages due to the statistical character of the received message and due to the nature of the channel will be considered.

The fundamental problem of communication is that of reproducing at one point or more points, or approximately, a message selected at another point. Frequently the messages have meaning; that is they refer to or are associated with certain physical configurations. These semantic aspects are irrelevant to the engineer and he may consider only the physical problem. The system he designs must be capable of carrying information without regard to what it carries or what it means. The system must be designed to operate for every possible selection, not just the one which will actually be chosen since this is unknown at the time of design.

If the number of messages in the set is finite then this number or any reasonable fraction of this number can be approached at a reasonable rate by a mechanism produced when the message is chosen. But if the set, all choices being equally likely, is infinite, then the number of possible states of the system is infinite. It follows that the number of possible messages is infinite. Although this definition may be physically unsatisfactory when we consider the audience of the system or the message and when we have a continuous range of messages, we will, in all cases we are reasonably legitimate assume

The logical measure measure is more convenient for various reasons:

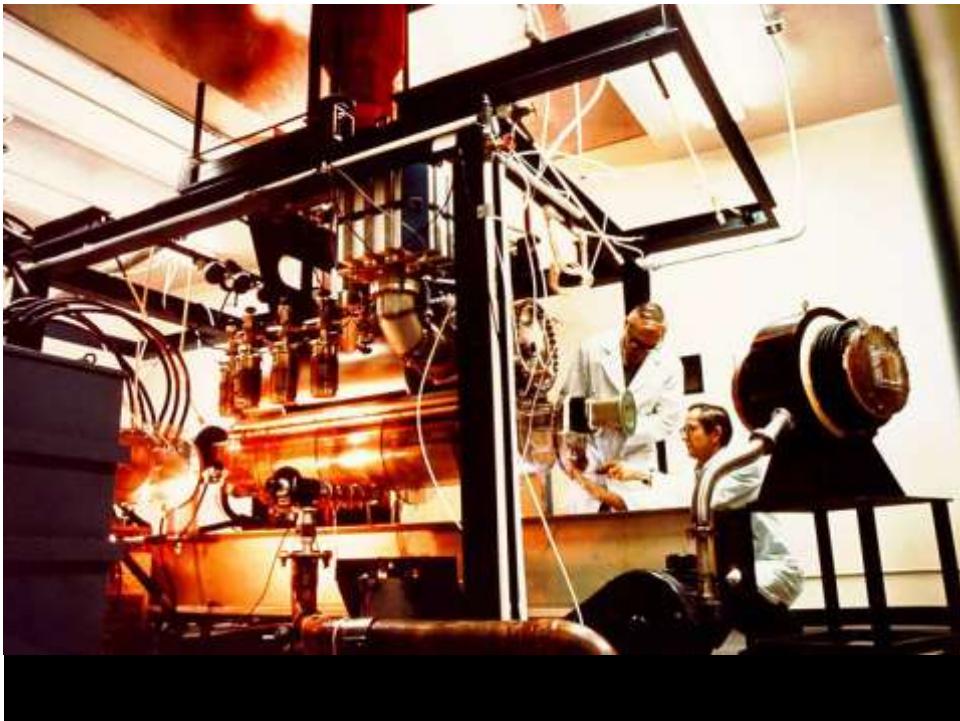
1. It is practically most useful. Possessing the properties required such as time, bandwidth, number of states, etc., as well as very closely with the logarithm of the number of possibilities. For example, two punched cards should have twice the capacity of one for information storage, and two identical channels twice the capacity of one for transmitting information.
2. It is a natural and intuitive feeling to the proper measure. This is closely related to 1) since for an arbitrary measure certain basic properties will always hold. One finds, for example, that two punched cards should have twice the capacity of one for information storage, and two identical channels twice the capacity of one for transmitting information.

The choice of a logarithmic base corresponds to the choice of a unit for measuring information. If the base 2 is used the resulting units may be called binary digits, or simply binary digits. A word suggested by J. W. Tukey³ is device reflecting middle positions, such as a relay or a flip-flop circuit, can store one bit of information. If each storage unit stores N bits, there will be N distinguishable possible states in 2^N , and $\log_2 N = N$. Hence base 10 is used the same may be called decimal digits. Since

$$\log_{10} M = \log_2 M / \log_2 10 = 3.321928 M$$

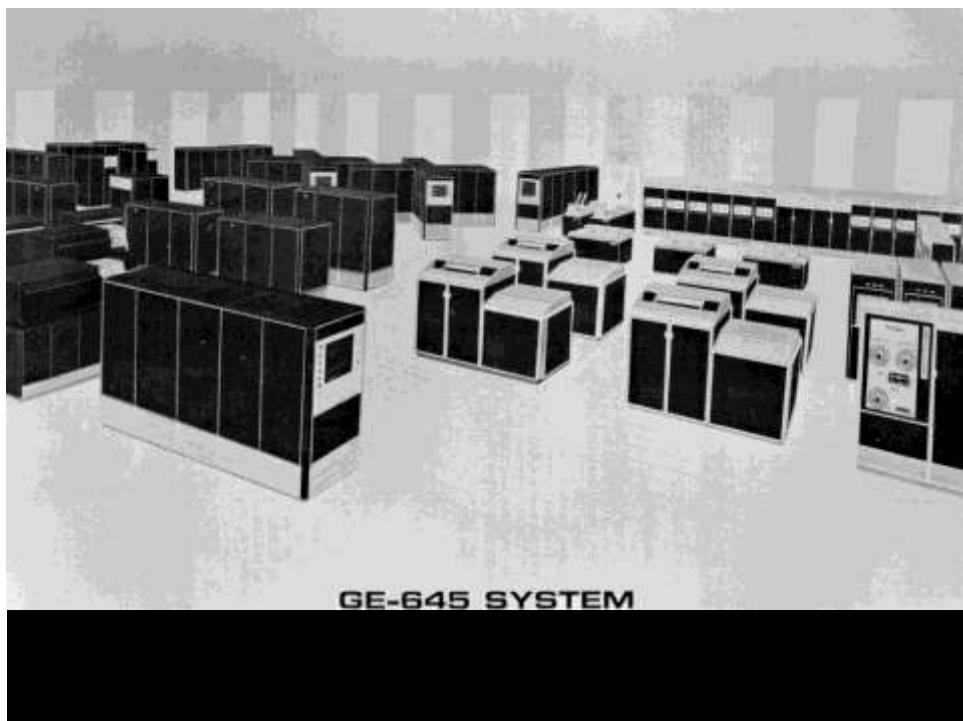
¹ Nyquist, H., "Certain Topics in Telegraph Transmission Theory," *Bell System Technical Journal*, April 1924, p. 326; "Certain Topics in Telecommunications Theory," *A.I.E.E. Trans.*, v. 47, April 1924, p. 617.

² Hartley, R. V. L., "Transmission of Information," *Bell System Technical Journal*, July 1928, p. 535.

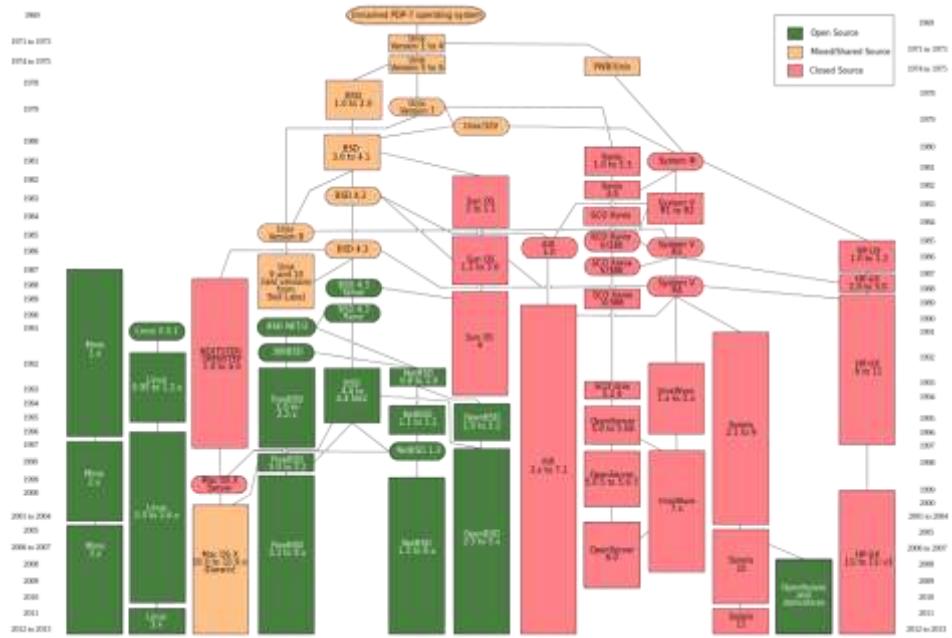












Why Unix is important

- Exemplar design
 - Technical contributions,
 - Impact
 - Development model
 - Widespread use
 - “unusual simplicity, power, and elegance”

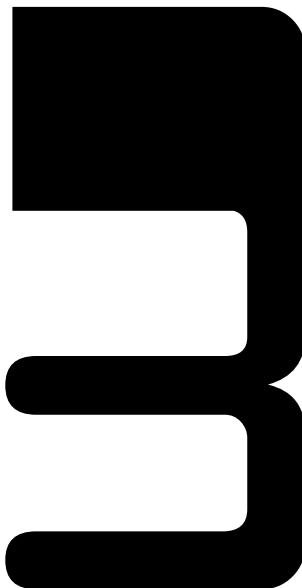


System technology

- Hierarchical file system
- Compatible file, device, networking, and inter-process I/O
- Pipes and filters architecture
- Virtual file systems
- The shell as a user-selectable regular process

Associated Technologies

- C and C++
- Parser and lexical analyzer generators
- Software development environments
- Document preparation tools and declarative markup
- Scripting languages
- TCP/IP networking
- Configuration management systems





A.T. & T. SETTLES ANTITRUST CASE; SHARES PATENTS

U.S. Hails Consent Decree as Major Victory—Company Calls Terms 'Stringent'

By ANTHONY LEWIS

Special to The New York Times
WASHINGTON, Jan. 21.—An antitrust suit against the American Telephone and Telegraph Company was settled today on terms described by Government lawyers as a major victory.

Herbert Brownell Jr., Attorney General, announced the signing of a consent decree in the Federal Court in Newark, N. J. Under the terms of the settlement A. T. & T. must:

• License 8,600 existing patents in all applicants without royalties.

• Licensen all its other patents, present and future, to any American concern at "reasonable and nondiscriminatory" rates.

• Get out of all business not directly connected with the communications field.

• Maintain uniform cost accounting methods for its manufacturing subsidiary, Western Electric.

One of 'Most Important'
Stanley N. Marcus, Assistant Attorney General in charge of the Justice Department's Antitrust Division, said the decree was "one of the most important" in antitrust history. Another department lawyer called it "miraculous."

In New York, Cleo F. Craig, president of A. T. & T., acknowledged that the terms of the consent decree were "stringent." However, he said, the settlement will leave intact "the unique combination and teamwork of the operating companies, the Bell Telephone Laboratories and the Western Electric Company that over the years has produced for the people of this country the finest, most widely used and most progressive telephone service in the world."

The A. T. & T. case was one of three major antitrust suits brought by the Government in the electronics field since World War II. The others, involving the Radio Corporation of America and International Business Machines, also are in negotiation for possible consent settlements. The I. B. M. negotiations are believed to be almost finished.

Through subsidiary Bell operating companies, A. T. & T. controls a majority of the country's telephone lines. Western Electric, its wholly owned subsidiary, makes the equipment for all Bell companies.

U. S. Pressed Civil Suit

On Dec. 21, 1954, the assets of A. T. & T. and the Bell system were estimated at \$13,000,000,000.

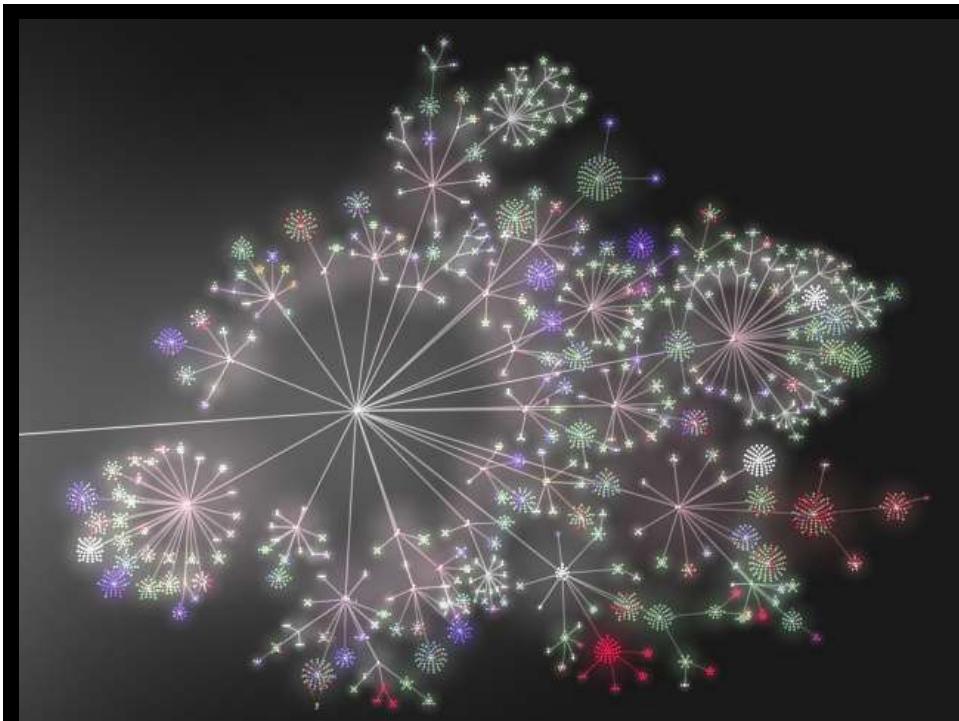
The Government complaint filed in 1949, charged that A. T. & T. and Western Electric had "unfairly restrained and monopolized trade and commerce in the manufacture, distribution, sale and installation of telephone equipment."

It was a civil suit. The Government was not calling for a fine but wanted the court to order changes in A. T. & T.'s structure. Specifically, the Government asked that the parent corporation give up its interest in Western Electric, that Western Electric be dissolved and its assets divided among three other companies.

The judgment entered today allows Western Electric to continue as manufacturer to the Bell System. However, several

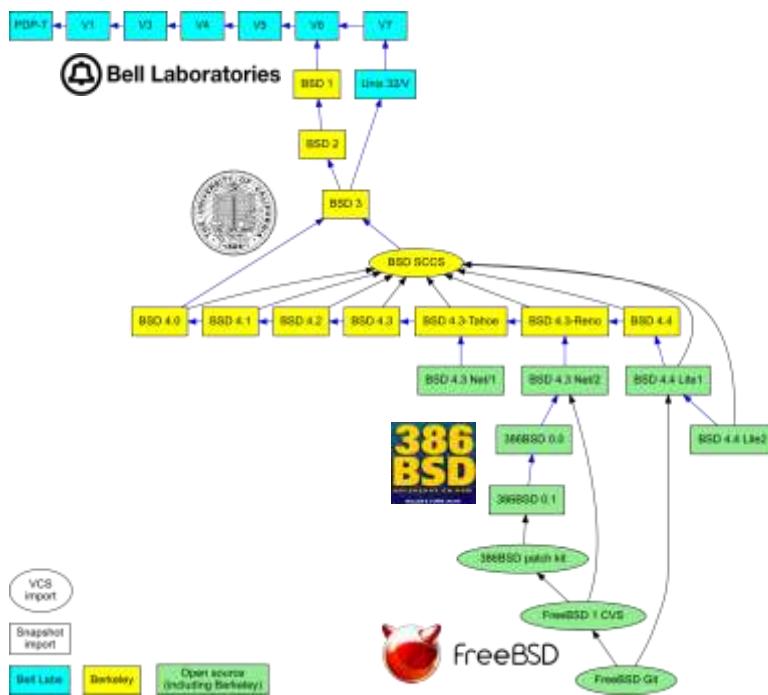
Continued on Page 15, Column 4

The New York Times
Published January 22, 1955
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Motivation

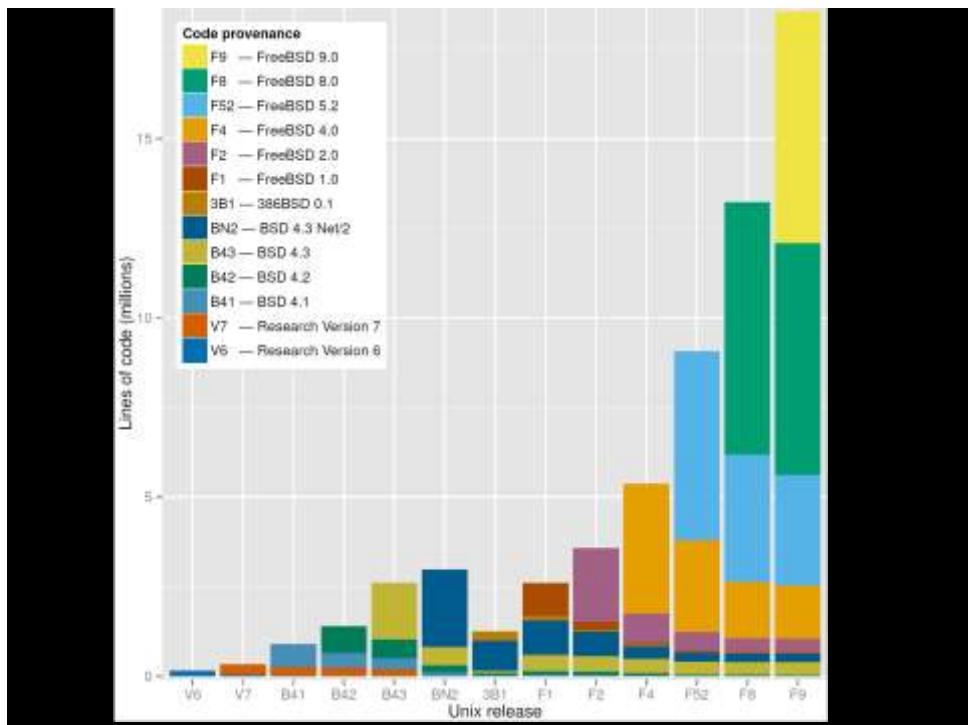
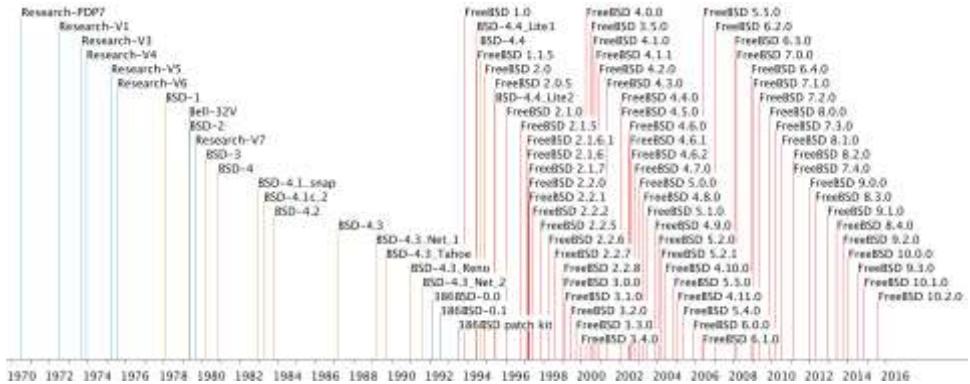
- Explore evolution of programming style
- Consolidate digital artifacts of historical importance
- Collect and record history that is fading away
- Provide a data set of digital archeology and repository mining

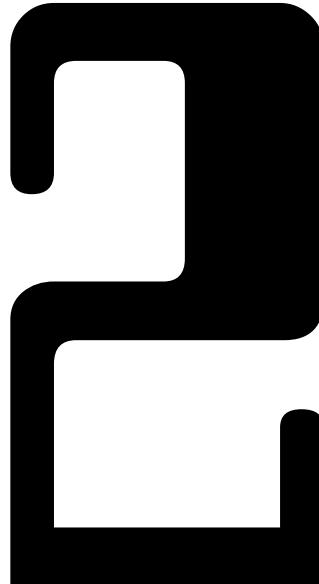


```
$ git checkout FreeBSD-RELEASE-10.0.0
$ git blame --on -c -c -t lib/libc/gen/timezone.c
usr/src/libc/gen/timezone.c (Dennis Ritchie 1979-01-10 14:58:45 -0500 76) static struct zone {
usr/src/libc/gen/timezone.c (Dennis Ritchie 1979-01-10 14:58:45 -0500 77)     int offset;
usr/src/libc/gen/timezone.c (Dennis Ritchie 1979-01-10 14:58:45 -0500 78)     char *stdzone;
usr/src/libc/gen/timezone.c (Dennis Ritchie 1979-01-10 14:58:45 -0500 79)     char *dstzone;
usr/src/libc/gen/timezone.c (Dennis Ritchie 1979-01-10 14:58:45 -0500 80) } zonetab[] = [
lib/libc/gen/timezone.c (Jordan K. Hubbard 1996-07-12 18:57:58 +0000 81)     {-1"60, "MET", "MET DST"),
[...]
lib/libc/gen/timezone.c (Jordan K. Hubbard 1996-07-12 18:57:58 +0000 96)     (-1)
use/src/lib/libc/gen/timezone.c (Bill Joy 1980-12-22 00:40:25 -0800 97) };
use/src/lib/libc/gen/timezone.c (Bill Joy 1980-12-22 00:40:25 -0800 98) ;
use/src/lib/libc/gen/timezone.c (Keith Bostic 1987-03-28 19:27:07 -0800 106) char *
lib/libc/gen/timezone.c (Ed Schouten 2009-12-05 19:31:38 +0000 107) _ctab(int zone, int dst)
lib/libc/gen/timezone.c (Rodney Grimes 1994-05-27 05:00:24 +0000 108) {
lib/libc/gen/timezone.c (David E. O'Brien 2002-02-01 01:08:48 +0000 109)     struct zone *zpt;
lib/libc/gen/timezone.c (David E. O'Brien 2002-02-01 01:08:48 +0000 110)     char sign;
use/src/lib/libc/gen/timezone.c (Bill Joy 1980-12-22 00:40:25 -0800 111) ;
use/src/lib/libc/gen/timezone.c (Keith Bostic 1987-03-28 19:27:07 -0800 112)     for (zp = zonetab; zp->offset != -1; ++zp)
use/src/lib/libc/gen/timezone.c (Keith Bostic 1987-03-28 19:27:07 -0800 113)         if (zp->offset == zone) {
use/src/lib/libc/gen/timezone.c (Dennis Ritchie 1979-01-10 14:58:45 -0500 114)             if (dst && zp->dstzone)
use/src/lib/libc/gen/timezone.c (Dennis Ritchie 1979-01-10 14:58:45 -0500 115)                 return(zp->dstzone);
use/src/lib/libc/gen/timezone.c (Dennis Ritchie 1979-01-10 14:58:45 -0500 116)             if (!dst && zp->stdzone)
use/src/lib/libc/gen/timezone.c (Dennis Ritchie 1979-01-10 14:58:45 -0500 117)                 return(zp->stdzone);
use/src/lib/libc/gen/timezone.c (Dennis Ritchie 1979-03-28 19:27:07 -0800 118)         }
use/src/lib/libc/gen/timezone.c (Keith Bostic 1987-03-28 19:27:07 -0800 119)     }
use/src/lib/libc/gen/timezone.c (Keith Bostic 1987-03-28 19:27:07 -0800 120)     if (zone < 0) {
use/src/lib/libc/gen/timezone.c (Bill Joy 1980-12-22 00:40:25 -0800 121)         zone = -zone;
use/src/lib/libc/gen/timezone.c (Keith Bostic 1987-03-28 19:27:07 -0800 122)         sign = '+';
use/src/lib/libc/gen/timezone.c (Keith Bostic 1987-03-28 19:27:07 -0800 123)     },
use/src/lib/libc/gen/timezone.c (Keith Bostic 1987-03-28 19:27:07 -0800 124)     else
use/src/lib/libc/gen/timezone.c (Keith Bostic 1987-03-28 19:27:07 -0800 125)         sign = '-';
lib/libc/gen/timezone.c (Warner Losh 1998-01-21 21:46:36 +0000 126)     (void)snprintf(czone,
sizeof(czone),
lib/libc/gen/timezone.c (Warner Losh 1998-01-21 21:46:36 +0000 127)         "%04c%0d;%02d", sign, zone, /
60, zone % 60);
lib/libc/gen/timezone.c (Rodney Grimes 1994-05-27 05:00:24 +0000 128)     return(czone);
lib/libc/gen/timezone.c (Rodney Grimes 1994-05-27 05:00:28 +0000 129) }
```

In Numbers ...

Metric	Unix history	Linux history
Start date	30/06/1970	17/09/1991
Start files	43	92
Start lines	11,500	917,812
End files	63,049	51,396
End lines	27,388,943	21,525,436
Data set size (.git)	1.1GB	1.0GB
Number of commits	495,622	611,735
Number of merges	2,523	48,821
Number of authors	973	18,465
Days with activity	13,004	5,126





dspinellis.github.io/unix-history-man

Evolution of Unix Facilities

1. User commands
2. System calls
3. C library functions
4. Devices and special files
5. File formats and conventions
6. Games et. al.
7. Miscellanea
8. System maintenance procedures and commands
9. System kernel interfaces



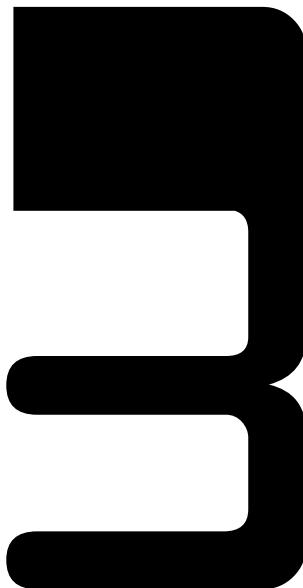
Evolution of Unix section 2: System calls

Facility name	Appearance Research V1	Research V1	Research V2	Research V3	Research V4	Research V5	Research V6	QSD 1	QSD 2	Get SDV	Research V7	QSD 3
umount	Research V1											
unlink	Research V1											
wall	Research V1											
wire	Research V1											
cd	Research V2											
dup	Research V2											
fdup	Research V2											
select	Research V2											
spmr	Research V2											
boot	Research V3											
cdev	Research V3											
dup2	Research V3											
tpre	Research V3											
rsave	Research V3											
pipe	Research V3											
times	Research V3											
readfd	Research V4											

Search in evolution tables

Disclaimers:

- The name of a facility may have been re-purposed over time.
- Facilities in sections 1, 6, 8 moved across sections over time. To allow a continuous view of their evolution, all have been re-located to the section of the most recent FreeBSD release. If they still existed at the time the evolution data of collected their names, reflect the evolution of the first stable code.







Numbers & Insights

Things to Take Away ...

- Architectural evolution in practice
 - 1.1GB Git repository
 - github.com/dspinellis/unix-history-repo
 - Open source project
 - github.com/dspinellis/unix-history-make
 - Architectural lessons to apply
-
- Joint work with **Paris Avgeriou** Professor (chair of Software Engineering) and Head of the SEARCH (Software Engineering and Architecture) Group of the Johann Bernoulli Institute for Mathematics and Computer Science at the University of *Groningen* (RuG).





PDP-7 [Unix] (1970)

- Kernel
- Layering and Partitioning
- System Call
- Data Scoping
- Interpreter
- Monolithic Implementation
- Process Management
- Descriptor Management
- Separation of File Metadata from File Naming
- Devices as Files
- File I/O
- Filesystem

```

between 0
    lne ema
    lac between i
    dae 9f+t
    lsz between
    laeq
    tad 9f+t i
    ema
    jmp 1f
    lac between i
    dae 9f+t
    lsz between
    laeq
    tad 9f+t i
    ema
    spa ema
    1f
    lsz between
    laeq
    ema
    jmp between i

copy 0
    =1
    tad copy i
    dae 8
    lsz copy
    =1

```

between 0
 dae 9f+t
 xct between i
 tad 9f+t lsz between
 spa
 jmp ff
 xct between i
 lsz between
 tad 9f+t
 spa ema
 i: lsz between. eme 9f+t
 jmp between i

Kernel

- 2584 lines
- Loads and executes user-level commands
- Provides the file abstraction
- Virtualizes the hardware interfaces
- Establishes ownership of files

Layering and Partitioning

adm.s	cat.s	dskio.s	init.s	s6.s
ald.s	check.s	dskres.s	lcase.b	s7.s
apr.s	chmod.s	dsksav.s	maksys.s	s8.s
as.s	chown.s	ds.s	s1.s	s9.s
bc.s	chrn.s	dsw.s	s2.s	scope.v
bi.s	cp.s	ed1.s	s3.s	sop.s
bl.s	db.s	ed2.s	s4.s	trysys.s
cas.s	dmabs.s	ind.b	s5.s	

Process Management (fork)

```

.fork!
jms lookfor 0 "not-used
    skip
    jms_error
    dae 9f4t
    lsr uniqpid
    lac uniqpid
    dae u.ac
    lav sysexit
    dae u.swapret
    lac 0200000
    tda u.ulistp i
    dae u.ulistp i
    jms dsksavp/ 07000
    lac 9f4t
    dae u.ulistp
    lac 0100000
    xor u.ulistp i
    dae u.ulistp i
    lac u.pid

```

Descriptor Management

```

fgeti 0
    jms between d01 d9
    jmp fget i
    clih mult 9
    lsz
    lsz

    tad ofilesp
    dac 9f+t
    dac .+2
    jms copyi .,; fnodes 3
    lsz fget
    jms fget i

    fputi 0
    lac 9f+t
    dac .+3
    jms copyi fnodes .,; 3
    jmp fput i
    t = t+1

```

Separation of File Metadata from File Naming

inodei i.flagst ,F,+1 i.nkpsl ,S,+7 i.uidl ,F,+1 i.nlksl ,S,+1 i.sizei ,S,+1 i.uniqi ,S,+1 , = inode+12	nameit 0 jms iget -1 tad namei i dac 9f+t+1 lsz namei lac i.flags and o20 sna jmp namei i -8 tad i.size cma irss 3 dac 9f+t sna jmp namei i dzm di
---	--

Devices as Files

```
ttyin:  
    <tt>;<yi>;<n 040;040040  
ttyout:  
    <tt>;<yo>;<ut>; 040040  
keybd:  
    <ke>;<yb>;<oa>;<rd>  
displ:  
    <di>;<sp>;<la>;<y 040  
sh:  
    <sh>; 040040;040040;040040  
system:  
    <sy>;<st>;<em>; 040040
```

File I/O

- open
- read
- write
- seek
- tell
- close

Filesystem

- creat
- rename
- link
- unlink

Interpreter

```

main $(
    extrn read, write;
    auto i, c, state, line 100;
) und.6

loop:
    state = i = 0;
loop1:
    c = read();
    if(c==0) return;
    if(c=='1' & state==0) state = 2;
    if((c<'0' & c>'9' & c<'a' & c>'z') & state==0) state = 1;
    line[i] = c;
    i = i+1;
    if(c!=012) goto loop1;
    if(state==2 & i==1) goto noi;
    write(' ');
    write(' ');
noi:
    i = 0;
loop3:
    c = line[i];
    write(c);
    i = i+1;
    if(c!=012) goto loop3;
    goto loop;
)

```

First Research Edition (Nov 1971)

- System Calls
- Binary-Code API
- Abstraction of Standard I/O
- Generic File I/O Layer
- User-Contributed Tools and Games
- The Shell as a User Program
- Interoperability through Documented File Formats
- Tree Directory Structure
- Mountable Filesystem Interface





Bell Laboratories

asym: Study of UNIX

date: September 14, 1972

from: T. R. Bashkow

Messrs. W. S. Bartlett
 D. P. Clayton
 D. H. Copp
 Mmes. G. J. Hansen
 J. Hints
 Mr. L. J. Kelly
 Miss R. L. Klein

Messrs. J. J. Ludwig
 J. F. Maranzano
 Mrs. G. Pettit
 Messrs. J. E. Ritacco
 B. A. Tague
 D. W. Vogel
 Mrs. L. S. Wright

On Tuesday, September 19, at 9:30 a.m. in
 Room 2A-418 at Murray Hill, I will give a talk on my study
 of the UNIX operating system. The emphasis will be on the
 structure, functional components, and internal operation
 of the system.

MH-8234-TRB-mbh

T. R. Bashkow

Copy to:
Mr. G. L. Baldwin*T.R.Bashkow*

```
/ initialize inodes for special files (inodes 1 to 40.)
      mov    $40,,r1 / set r1=i-node-number 40.
1:     jsr    r0,iget / read i-node 'r1' from disk into inode area of
                  / core and write modified inode out (if any)
      mov    $100017,i.flags / set flags in core image of inode to indi-
                  / cate allocated, read (owner, non-owner),
                  / write (owner, non-owner)
      movb   $1,i.nlks / set no. of links = 1
      movb   $1,i.uid / set user id of owner = 1
      jsr    r0,setimod / set imod=1 to indicate i-node modified, also
                  / stuff time of modification into i-node
      dec    r1 / next i-node no. = present i-node no.-1
      bgt    1b / has i-node 1 been initialized; no, branch

/ initialize i-nodes r1.....,47. and write the root device, binary, etc.,
/ directories onto fixed head disk. user temporary, initialization prog.
```

Issue D Date 3/17/72 ID IMO.1-1 Section E.0 Page 4

UNIX IMPLEMENTATION

```

/* u0 == unix
cold = 0
orig = 0 /* orig = 0, relocatable
phda = 177403 / disk address reg
dr00 = 177404 / disk status reg
dr02 = 177404 / control status reg
year = 177400 / receiver status reg
dr04 = 177402 / receiver buffer reg
dr05 = 177403 / transmitter status reg
dr06 = 177404 / transmitter buffer reg
dr07 = 177405 / zero buffer reg
test = 177440 / disk tape control status
tcr0 = 177342 / disk tape command reg
tcr1 = 177343 / disk tape control reg
tcr2 = 177344 / bus addres
tcr3 = 177346 / data reg
dcm = 177450 / disk control status
dcm = 177450 / error status session
lks = 177548 / clock status reg
prv = 177552 / reportage reader status
prv = 177553 / punch buffer
prv = 177554 / punch buffer
prv = 177556 / punch buffer
prv = 177557 / processor status
halt = 0
wait = 1
rti = 2
nproc = 16 / number of processes
nfiles = 50;
ntry = 84;
nbuf = 1
/* if cold = 1, then 00000000 / ignored if cold = 0
nbuf = 2
.cold
core = core+40000 / specifies beginning of user's core
core = core+10000 / specifies end of user's core (4096 w)
.cold
#1) d11 init by copy
#2) fcopy2 = idle in tr
#3) unknown / trace and trap (see Sec. 8.1 page 1)
#4) unknown / trap
#5) unknown / trap
#6) unknown / trap
#7) unknown / set
#8) unknown / sys
#9) unknown / sys

```

Issue D Date 3/17/72 IP IMD.1-1 Section E.0

UNIX IMPLEMENTATION

```

mov $idata, u0 / r0 base addr. of assembled directories.
mov $uoff,u0,top / pointer to u.off in u.ofop (holds file
                   / offset)

14 mov ($r0)=u0 / "0" in the assembled directory
      / header signals last
beq r0 / assembled directory has been written onto drum
jst r0,loop / set the bit indicating bit for i-node 'r'
bsh $u0,(r2) / set the bit to indicate the i-node is not
               / available
jst r0,writer / read inode 'r' from disk into inode area of
               / core and then write it onto drum (if any)
mov ($r0)=u0,top / assembled directory header
movb ($r0)=u0,nlink / set user id of owner from header
movb ($r0)=u0,inode / set index to which inode modified; also,
                     / current time to modification into i-node
mov ($r0)=u0,ccount / set byte count for write call equal to
                     / size of directory
mov $r0,u_base / set buffer address for write to top of directory
clr $r0,u_base / clear file offset in r0 to 0
jst r0,writes / read drum superblock
15 tsetb abd+1 / has I/O request been honored (for drum)?
bne 15, no, continue in idle.

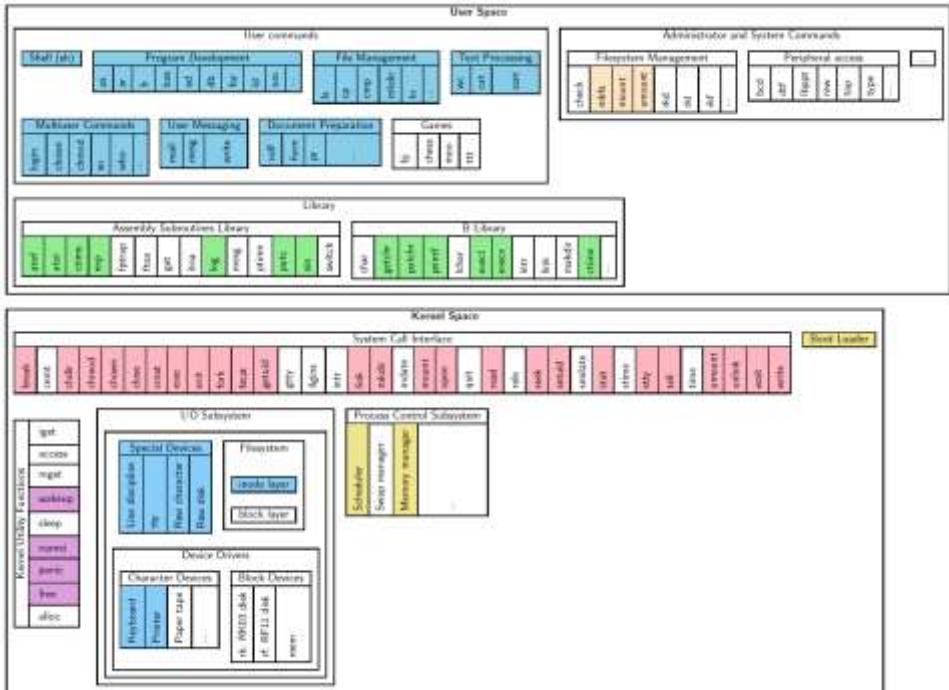
16 decb sysflag / normally sysflag=0, indicates executing in system
sys exec; zf; tf / generated trap interrupt; trap vector =
                   / sysent; 0
br panic / execute file/etc/init

17 2f:0
2f:0 _____ This is file 2f:0 L1 in E1.4 See 6.10
Q/etc/init:0 / UNIX looks for strings term, noted by null\0

panic:
clr pc
18 dec 40
inc 18
dec 40
inc 18
jmp *$13700 / rom loader address

```

Issue D Date 3/17/72 IP IMD.1-1 Section E.0 Page 5



First Edition Unix 1972 FreeBSD 11.1 2018

sysrele	/ 0	0 { int nosys(void); } syscall nosys_args int
sysexit	/ 1	1 { void sys_exit(int rval); } exit \
		sys_exit_args void
sysfork	/ 2	2 { int fork(void); }
sysread	/ 3	3 { ssize_t read(int fd, void *buf, \
		size_t nbytes); }
syswrite	/ 4	4 { ssize_t write(int fd, const void *buf, \
sysopen	/ 5	size_t nbytes); }
sysclose	/ 6	5 { int open(char *path, int flags, int mode); }
syswait	/ 7	6 { int close(int fd); }
syscreat	/ 8	7 { int wait4(int pid, int *status, \
		int options, struct rusage *rusage); }
syslink	/ 9	8 { int creat(char *path, int mode); }
sysunlink	/ 10	9 { int link(char *path, char *link); }
		10 { int unlink(char *path); }

Issue D Date 3/17/72

ID IMO.1-1-

Section E.1

Page 1

The Shell as a User Program

11/3/71

PASSWD (V)

NAME	passwd -- password file
SYNOPSIS	--
DESCRIPTION	<u>passwd</u> contains for each user the following information: name (login name) password numerical user ID default working directory <u>program to use as Shell</u>

This is an ASCII file. Each field within each user's entry is separated from the next by a colon. Each user is separated from the next by a new-line. If the password field is null, no password is demanded; if the Shell field is null, the Shell itself is used.

Abstraction of Standard I/O

11/3/71

SH (I)

Two characters cause the immediately following string to be interpreted as a special argument to the shell itself, not passed to the command. An argument of the form "<arg" causes the file arg to be used as the standard input file of the given command; an argument of the form ">arg" causes file "arg" to be used as the standard output file for the given command.

Interoperability through Documented File Formats

V. FILE FORMATS

a.out	assembler and loader output
archive	archive file
bppt	binary paper tape format
core	core image file
directory	directory format
file system	file system format
passwd	password file
uids	map names to user ID's
utmp	logged-in user information

Format	Description	Clients
a.out	Assembler and linker output	as, ld, strip, nm, un
Archive	Object code libraries	ar, ld
Core	Crashed program image	Kernel, db
Directory	File system directories	du, find, ls, ln, mkdir, rmdir
File system	File system format	check, dump,* mkfs, restor*
Ident	GECOD ident card format	opr
Password	User accounts and passwords	chown, find, getpw,* login,* ls, passwd*
Tape*	DECtape file format	mt,* tap*
Uid	User identifier to name map	chown
utmp	Logged in users	init, login,* who,* write*
wtmp*	Users login history	acct, date, init, login, tacct, who

User-Contributed Tools and Games

VI. USER MAINTAINED PROGRAMS

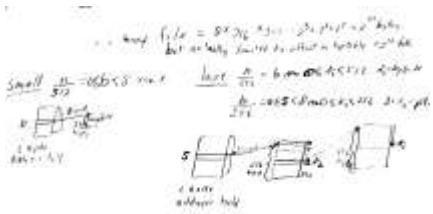
```

basic ..... DEC supplied BASIC
bj ..... the game of black jack
cal ..... print calendar
chess ..... the game of chess
das ..... disassembler
dli ..... load DEC binary paper tapes
dpt ..... read DEC ASCII paper tapes
moo ..... the game of MOO
sort ..... sort a file
ttt ..... the game of tic-tac-toe

```

Tree Directory Structure

- mkdir(II)
- chdir(II)
- chdir(I)
- find(I)
- ln(I)
- ls(I)
- stat(I)
- mkdir(I)
- mv(I)
- rm(I)
- rmdir(I)



Mountable Filesystem Interface

- mount(II)
- umount(II)
- mount(I)
- umount(I)

Second Research Edition (Jun 1972)

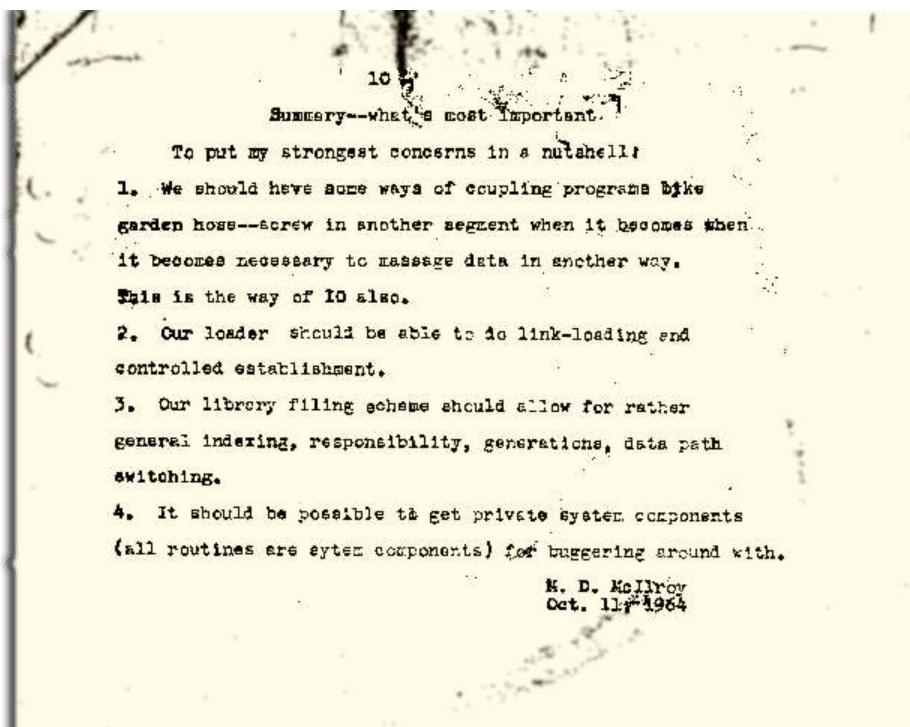
- Software Library

```
III. SUBROUTINES
stan ..... arctangent
atof ..... convert ASCII to floating
atoi ..... convert ASCII to integer
stoi ..... floating-point constants
const ..... convert time to ASCII
ctime ..... exponential function
exp ..... floating-point simulator
ftrap ..... convert floating to ASCII
ftoa ..... communicate with GCOS
gerts ..... get character
getc ..... compute hypotenuse
hypot ..... convert integer to ASCII
itoa ..... logarithm base e
log ..... print string on typewriter
mesg ..... read name list
nlist ..... print time
ptime ..... write character or word
putc ..... quicker sort
qsort ..... storage allocator
salloc ..... sine, cosine
sin ..... square root
sqrt ..... transfer depending on value
switch .....
```



Third Research Edition (Feb 1973)

- Pipes and Filters



Fourth Research Edition (Nov 1973)

- Structured Programming
- User Groups
- Language-Independent API
- Data Structure Definition Reuse
- Dynamic Resource Management
- Device Driver Abstraction
- Buffer Cache

Structured Programming

- Kernel implemented in “New B”
- 6373 lines New B
- 768 lines PDP-11 assembly
- 105 functions + 50 assembly symbols
- vs 248 global symbols in the First Edition

Language-Independent API

PIPE (II)

8/5/73

PIPE (II)

NAME

pipe — create a pipe

SYNOPSIS

```
(pipe = 42.)
sys pipe
(read file descriptor in r0)
(write file descriptor in r1)
pipe(fildes)
int fildes[2];
```

DESCRIPTION

The *pipe* system call creates an I/O mechanism called a pipe. The file descriptors returned can be used in read and write operations. When the pipe is written using the descriptor returned in r1 (resp. fildes[1]), up to 4096 bytes of data are buffered before the writing process is suspended. A read using the descriptor returned in r0 (resp. fildes[0]) will pick up the data.

It is assumed that after the pipe has been set up, two (or more) cooperating processes (created by subsequent *fork* calls) will pass data through the pipe with *read* and *write* calls.

The shell has a syntax to set up a linear array of processes connected by pipes.

Read calls on an empty pipe (no buffered data) with only one end (all write file descriptors closed) return an end-of-file. *Write* calls under similar conditions are ignored.

Data Structure Definition & Reuse

```
buf.h    filsys.h  proc.h   text.h  
conf.h   inode.h   reg.h    tty.h  
file.h   param.h  systm.h  user.h
```

Dynamic Resource Management

```
int      coremap[CMAPSIZ];  
int      swapmap[SMAPSIZ];  
  
struct map {  
    char *m_size;  
    char *m_addr;  
};  
  
malloc(mp, size)  
struct map *mp;  
{  
    ...  
}
```

Device Driver Abstraction

IV. SPECIAL FILES

cat	phototypesetter interface
da	voice response unit
dc	DC-11 communications interface
dn	dn11 ACU interface
dp	dp11 201 data-phone interface
kl	KL-11/TTY-33 console typewriter
mem	core memory
pc	PC-11 paper tape reader/punch
rf	RF11/RS11 fixed-head disk file
rk	RK-11/RK03 (or RK05) disk
rp	RP-11/RP03 moving-head disk
tc	TC-11/TU56 DECTape
tiu	Spider interface
tm	TM-11/TU-10 magtape interface
vs	voice synthesizer interface
vt	11/20 (vt01) interface

Driver Interface

```

struct {
    int (*d_open)();
    int (*d_close)();
    int (*d_strategy)();
} bdevsw[];

struct {
    int (*d_open)();
    int (*d_close)();
    int (*d_read)();
    int (*d_write)();
    int (*d_sgtty)();
} cdevsw[];

```

Buffer Cache

Fourth Edition

```
#define B_READ  01
#define B_DONE   02
#define B_ERROR  04
#define B_BUSY   010
#define B_XMEM   060
#define B_WANTED 0100
#define B_RELOC   0200
#define B_ASYNC   0400
#define B_DELWRI  01000
```

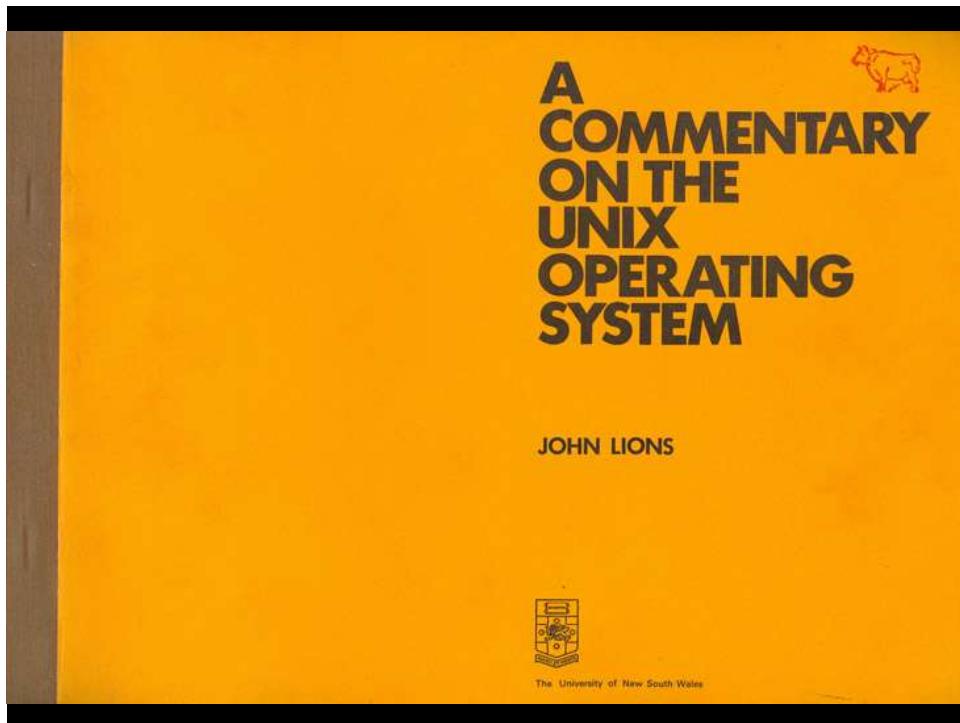
FreeBSD 11.1

```
#define B_ASYNC 0x00000004
/* Start I/O, do not wait.
 */
[...]
#define B_DELWRI 0x00000080
/* Delay I/O until buffer
reused. */
#define B_DONE 0x00000200
/* I/O completed. */
```

Fifth Research Edition (Jun 1974)

- Command Files

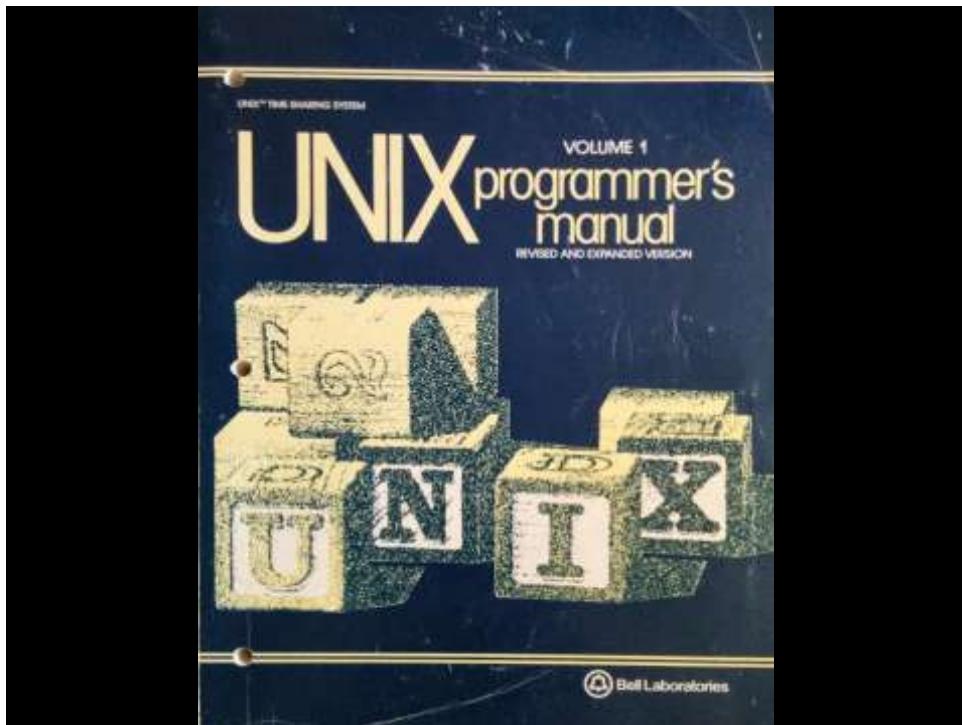
```
chdir /usr/source/s3
cc -c ctime.c
ar r /lib/liba.a ctime.o
rm ctime.o
chdir /usr/source/s1
cc -s -n date.c
cp a.out /bin/date
cc -s -n dump.c
cp a.out /bin/dump
cc -s -n ls.c
cp a.out /bin/ls
rm a.out
```



Sixth Research Edition (May 1975)

- Portable C Library

alloc.c	clenf.c	makbuf.c	scan1.c
calloc.c	copen.c	maktab.c	scan2.c
cclose.c	cputc.c	nexch.c	scan3.c
ceof.c	cwrd.c	nodig.c	system.c
cerror.c	dummy.s	printf.c	tmpnam.c
cexit.c	ftoa.c	putch.c	unget.c
cflush.c	getch.c	puts.c	unprnt.s
cfree.c	gets.c	relvec.c	wdleng.c
cgetc.c	getvec.c	revput.c	
ciodec.c	iehzap.c	run	



Seventh Research Edition (Jan 1979)

- Unix as a Virtual Machine
- Dynamic Memory Allocation
- Static Analysis
- Environment Variables
- Language Development Tools
- Domain-Specific Languages
- Filesystem Directory Hierarchy

Unix as a Virtual Machine

Also, about this time [1973] I had a fateful discussion with Dennis, in which he said

“I think it may be easier to port Unix to a new piece of hardware than to port a complex application from Unix to a new OS”

— Steve Johnson

Dynamic User Memory Allocation

- malloc(3), free(3)
- 26 programs: awk cc col cron dc dcheck diff ed eqn expr graph icheck learn ls m4 neqn nm quot ratfor spline struct tar tsort uucp xsend quiz
- stdio(3), mp(3)

Static Analysis



Environment Variables

- KEY=value
- Kernel
- Shell
- C Library

ENVIRON(5) UNIX Programmer's Manual ENVIRON(5)

NAME environ — user environment

SYNOPSIS

```
extern char **environ;
```

DESCRIPTION

An array of strings called the "environment" is made available by `environ()` when a process begins. By convention these strings have the form "name=value". The following names are used by various commands:

PATH The sequence of directory prefixes that `sh`, `zsh`, `xterm()`, etc., apply in searching for a file known by an incomplete path name. The prefixes are separated by ':'. `login()` sets `PATH=/bin:/usr/bin`.

HOME A user's login directory, set by `login()` from the password file `passwd(5)`.

TERM The kind of terminal for which output is to be prepared. This information is used by commands, such as `vi(1)` or `ps(1)`, which may exploit special terminal capabilities. See `term(7)` for a list of terminal types.

Further names may be placed in the environment by the `envir` command and "name=value" arguments to `sh(1)`, or by `env(2)`. It is advised to conflict with certain Shell variables that are frequently exported by "profile" files: `MAIL`, `PS1`, `PS2`, `IFS`.

SEE ALSO

`exec(2)`, `sh(1)`, `term(7)`, `login(1)`

Language Development Tools

- lex(1)
- yacc(1)
- 12 clients: awk bc
cpp egrep eqn lex
m4 make pcc
neqn struct
-



Domain-Specific Languages

- sh
- awk
- sed
- find
- expr
- egrep
- m4
- make

Filesystem Directory Hierarchy



First and Second Berkeley Software Distributions (1978)

- Software Packages
 - csh
 - ex
 - Mail
 - Pascal
 - termlib

3BSD (1979)

- Virtual Memory Paging
 - vm_*.c
 - 2808 out of 16039 C source code
 - 17% of kernel source code
 - vread(2), vwrite(2), vfork(2)



4BSD (Oct 1980)

- Regular Expression Library: regex(3)
- Optimized Screen Handling

Regular Expression Library: regex(3)

- 5 implementations: awk, sed, ed, grep, expr
- 1 client: more(1)
- 2 more by 4.3: dbx(1), rdist(1)
- 4 replacements in FreeBSD: ed, grep, sed, expr

Optimized Screen Handling

- curses(3)
- termcap(5)



4.2BSD (Sep 1983)

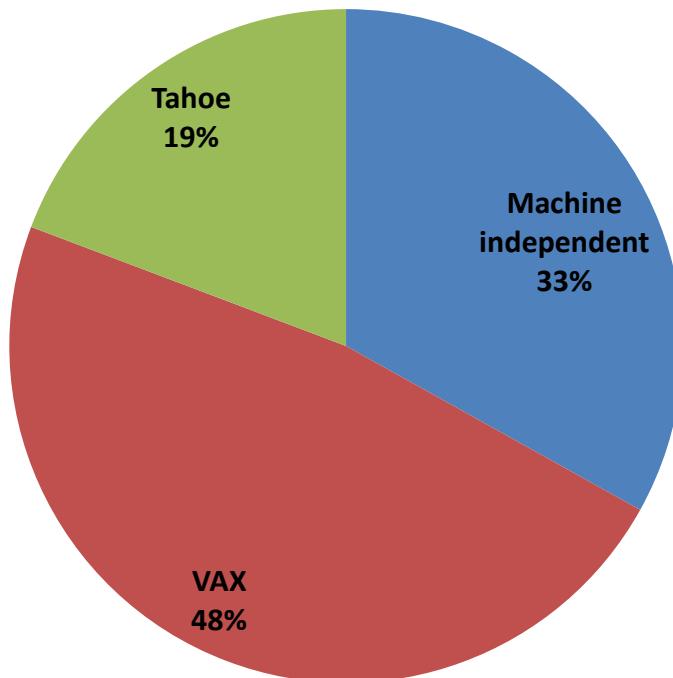
- Internet Protocol Family
 - ARP, IP, TCP, UDP, ICMP
- Local and Remote Interprocess Communication
 - socket(2), etc.
- Network and User Database Access
 - getfsent(3x), getgrent(3), gethostent(3n),
getnetent(3n), getprotoent(3n), getpwent(3),
getservent(3n)
- Pseudo-Terminal Driver
 - pty(4)

- 2 library functions
 - rcmd(3x)
 - rexec(3x)
- 11 system daemons
 - comsat(8c)
 - ftpd(8c)
 - gettable(8c)
 - implogd(8c)
 - rexecd(8c)
 - rlogind(8c)
 - af(8c)
 - rshd(8c)
 - rwhod(8c)
 - telnetd(8c)
 - tftpd(8c)
- 8 user-mode programs
 - ftp(1c)
 - rlogin(1c)
 - rsh(1c)
 - talk(1c)
 - telnet(1c)
 - tftp(1c)
 - whois(1c)
 - sendmail(1c)

System call	Uses
bind	23
connect	15
accept	13
select	12
listen	11
sendto	10
shutdown	9
recvfrom	8
getsockname	6
recv	2
send	2
sendmsg	1
getsockopt	0
recvmsg	0
socketpair	0

4.3BSD Tahoe (Jun 1988)

- Multiple CPU Architecture Support
 - VAX
 - CCI Power 6/32 (Tahoe)
- Timezone Handling
 - Community contribution
 - Separation of timezone rules from code



4.3BSD Reno (Jun 1990)

- Kernel Packet Forwarding Database
 - route(4)
 - routed(8), XNSrouted(8)
- Virtual Filesystem Interface
 - ...

vnode

```
/*
 * Operations on vnodes.
 */
struct vnodeops {
    int     (*vn_lookup)(            /* ndp */;
    int     (*vn_create)(           /* ndp, fflags, vap, cred */ );
    int     (*vn_mknod)(           /* ndp, vap, cred */ );
    int     (*vn_open)(             /* vp, fflags, cred */ );
    int     (*vn_close)(            /* vp, fflags, cred */ );
    int     (*vn_access)(           /* vp, fflags, cred */ );
    int     (*vn_getattr)(          /* vp, vap, cred */ );
    int     (*vn_setattr)(          /* vp, vap, cred */ );

    int     (*vn_read)(              /* vp, uiop, offp, ioflag, cred */ );
    int     (*vn_write)(             /* vp, uiop, offp, ioflag, cred */ );
    int     (*vn_ioctl)(             /* vp, com, data, fflag, cred */ );
    int     (*vn_select)(            /* vp, which, cred */ );
    int     (*vn_mmap)(              /* vp, ..., cred */ );
    int     (*vn_fsync)(             /* vp, fflags, cred */ );
    int     (*vn_seek)(              /* vp, (old)offp, off, whence */ );
};

...
```

4.3BSD Net/2 (Jun 1991)

- Stream I/O Functions
 - `funopen(3)`
 - GNU `funopencookie(3)` added in FreeBSD 11

4.4BSD (Jun 1994)

- Stackable Filesystems
 - `mount_null(8)`
 - `mount_union(8)`
- Generic System Control Interface (MIB)
 - `sysctl(1)`
 - `sysctl(3)`
 - `sysctl(9)`



386BSD Patch Kit (1992-1993)

- Organized Community Contributions
 - From open source software ...
 - ... to an open source **project**
- Patch metadata
 - title
 - author
 - description
 - prerequisites



FreeBSD 1.1 (May 1994)

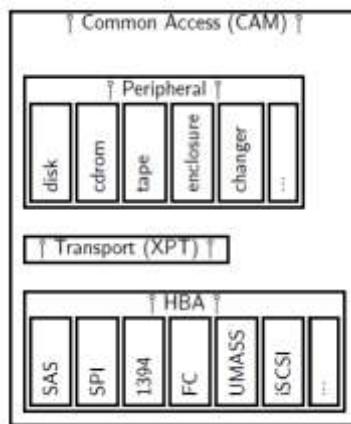
- Package Manager
 - Patch
 - Compile
 - Install
 - Uninstall
 - Handling of dependencies

FreeBSD 2.0 (Nov 1994)

- Process Filesystem
 - procfs(5)
- Dynamically Loadable Kernel Modules
 - lkm(4), then kld(4)
 - device drivers
 - file systems
 - emulators
 - system calls
 - **992** modules in FreeBSD 11.1

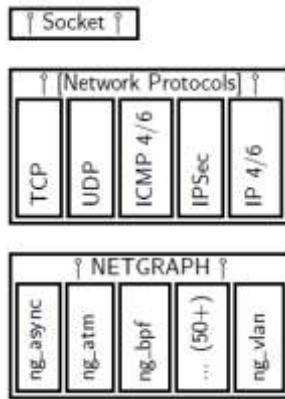
FreeBSD 3.0.0 (Jan 1999)

- Common Access Method I/O Subsystem (CAM)



FreeBSD 3.4.0 (Dec 1999)

- Graph-based Kernel Networking and User Library (NETGRAPH)

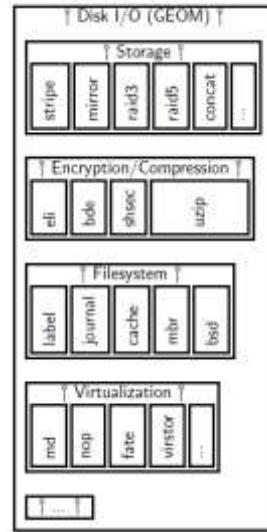


FreeBSD 4.0.0 (Mar 2000)

- OpenSSL Secure Sockets Layer and Transport Layer Security framework
 - Version 0.9.4
 - 1127 files, 227118 lines
 - libssl(3), libcrypto(3), openssl(1)
- Jail: Isolate a process and its descendants

FreeBSD 5.0.0 (May 2006)

- Modular Disk I/O Request Transformation Framework (GEOM)



FreeBSD 5.3.0 (Nov 2004)

- Streaming Archive Access Library
- Miniport Driver Wrapper

FreeBSD 7.0.0 (Feb 2008)

- ZFS Filesystems and Storage Pools

FreeBSD 7.1.0 (Dec 2008)

- Dynamic Tracing

FreeBSD 8.0.0 (Nov 2009)

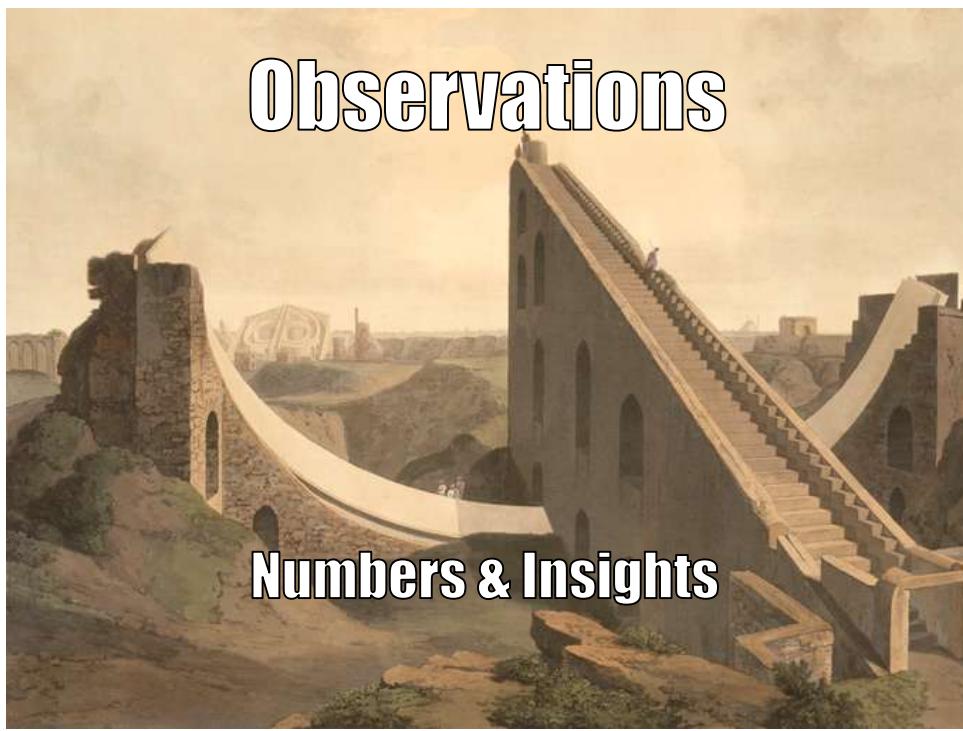
- Packet Capture Library
 - pcap(3)

FreeBSD 9.0.0 (Jan 2012)

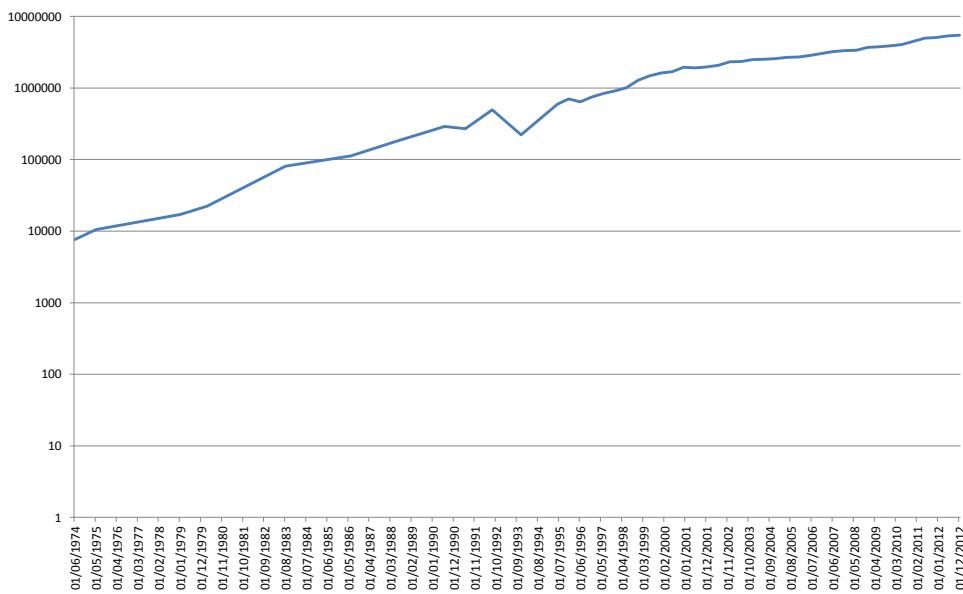
- Infiniband / RDMA High-Speed Low-Latency Switched-Fabric Interconnect Library

Observations

Numbers & Insights



Unix Kernel (Research, BSD, FreeBSD) LOC

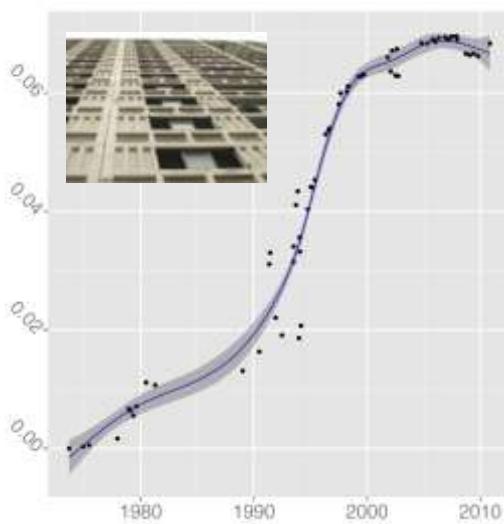




Modularity increases with code size

Increase in number of
static declarations /
statement

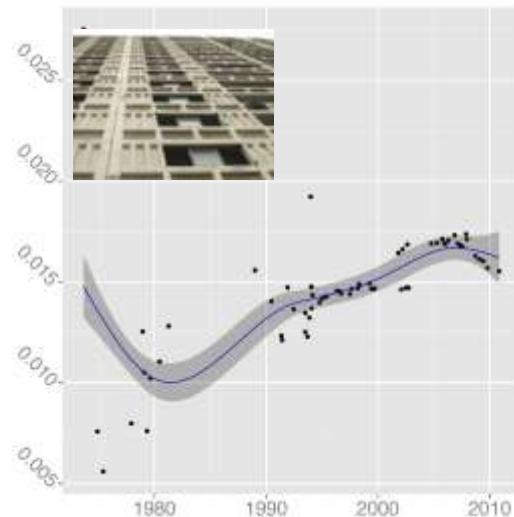
static short splice;



Modularity increases with code size

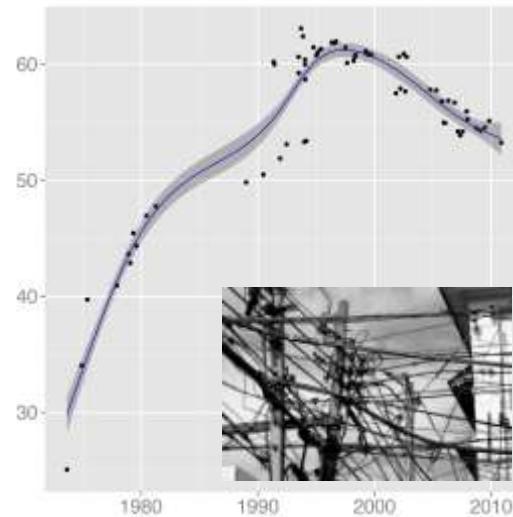
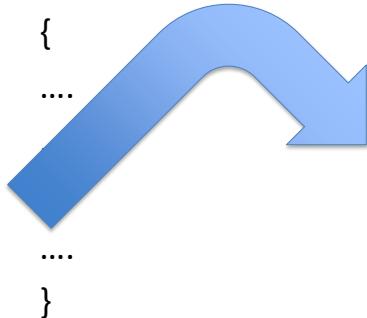
Increase in number of
#include directives /
line

#include "if_uba.h"



Software complexity evolution follows self correction

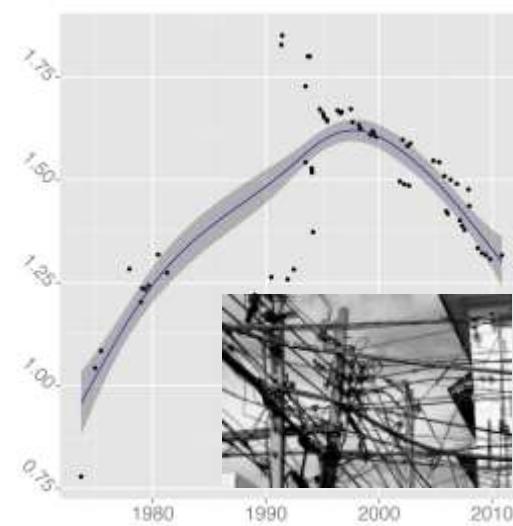
Mean lines / function



Software complexity evolution follows self correction

Mean statement
nesting

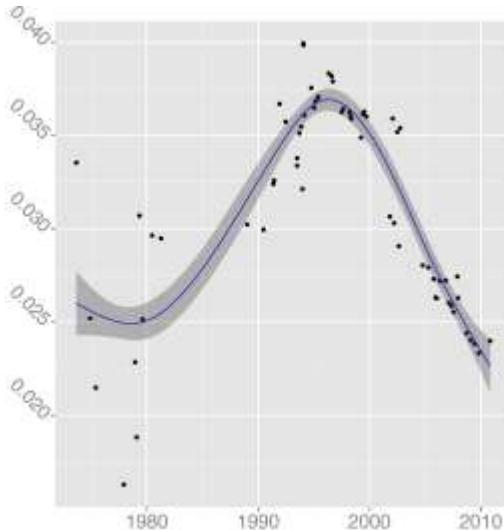
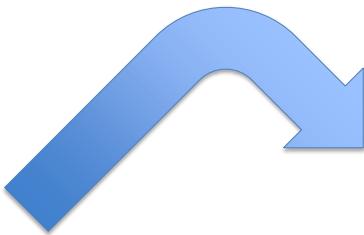
if (a)
while (b)
for (;;)
if (c)



Software complexity evolution follows self correction

Density of C
preprocessor non-
include directives

#define
#if



Letters to the Editor

Go To Statement Considered Harmful

Key Words and Phrases: go to statement, jump statements, branch statements, conditional clause, alternative clause, repetitive clauses, program readability, program reusability, C99 Committee: 0.22, 0.23, 0.24

Eduardo

For a number of years I have been familiar with the observation that the quality of programming is a decreasing function of the density of go-to statements in the programs they produce. More recently I discovered why the use of go-to statements has such disastrous effects, and I became convinced that things to statements should be abolished from all "higher level" programming languages (i.e., everything except, perhaps, plain machine code). As a result of this conviction, after two months of silence in the discussion, I now submit my resignation for publication because no recent discussion in which the subject turned up, I have been urged to do so.

My first remark is that, although the programmer's activity ends when he has constructed a correct program, the process taking place inside several of his programs is the true subject matter here. For it is this process that is responsible for the desired effect (i.e., to cause programs that in principle cannot fail to satisfy the detailed specifications. Yet, once the program has been made, the "walking" of the corresponding process is guided by the variables).

My second remark is that our intellectual powers are rather

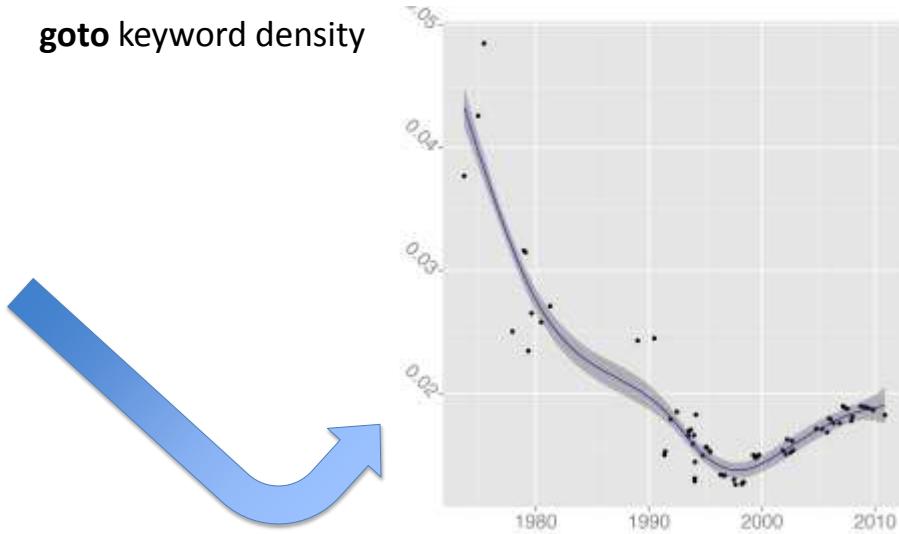
dynamically programs is only characterized when we also give to which cell of the procedure we enter. With the inclusion of procedures we can characterize the state of the process via a sequence of several indices, the length of this sequence being equal to the dynamic depth of procedure calling.

Let us now consider repetition clauses (i.e., while B repeat A or repeat A until B). Logically speaking, such clauses are now superfluous, because we can express repetition with the aid of recursive procedures. For example, if I want to enter a while clause on the one hand, repetition clauses are not needed; instead, we can proceed quite comfortably with present day facts equipped, on the other hand, the recursive pattern known as "indication"; makes us well equipped to start our intellectual grasp on the processes generated by repetition clauses. With the inclusion of indication clauses terminal indices are no longer sufficient to describe the dynamic state of the process. We need to introduce a repetition clause, however, we can associate a so-called "dynamic index," internally bearing the orbital number of the corresponding current repetition. An repetition clauses (not an procedure call) may be applied similarly, we find that now the progress of the process can always be uniquely characterized by a (finite) sequence of indices and dynamic indices.

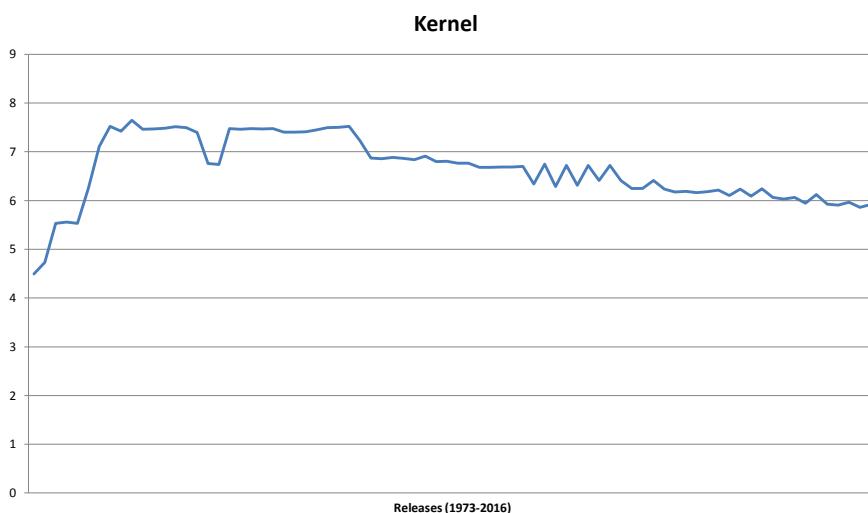
The main point is that the use of these indices are outside programmer's control; they are generated, either by the writer-up of his program or by the dynamic evolution of the process whether he wishes or not. They provide independent coordinates in which to describe the progress of the process.

EDSGER W. DIJKSTRA
*Technological University
Eindhoven, The Netherlands*

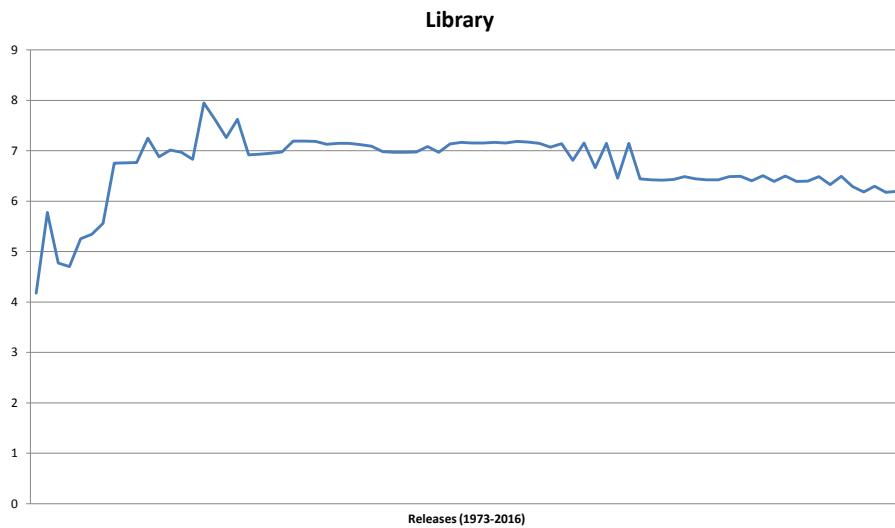
Software complexity evolution follows self correction



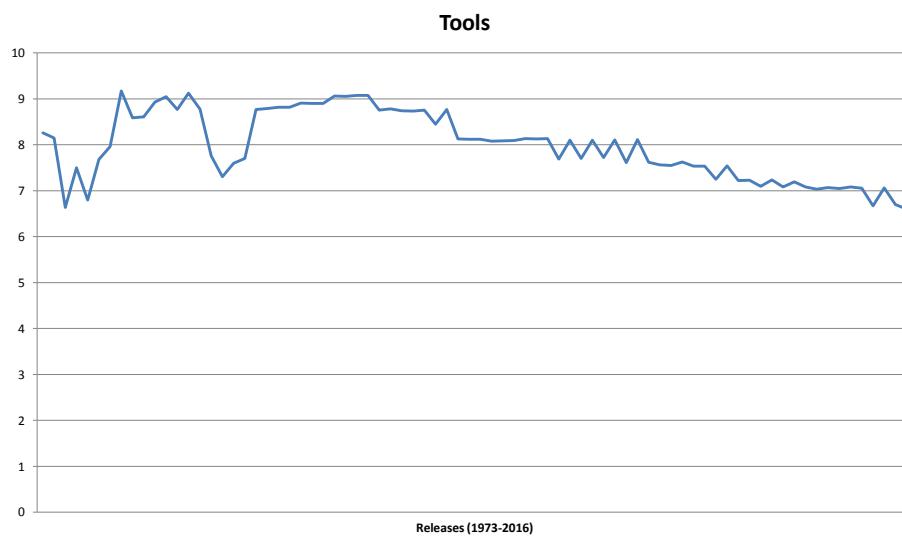
Cyclomatic Complexity



Cyclomatic Complexity



Cyclomatic Complexity

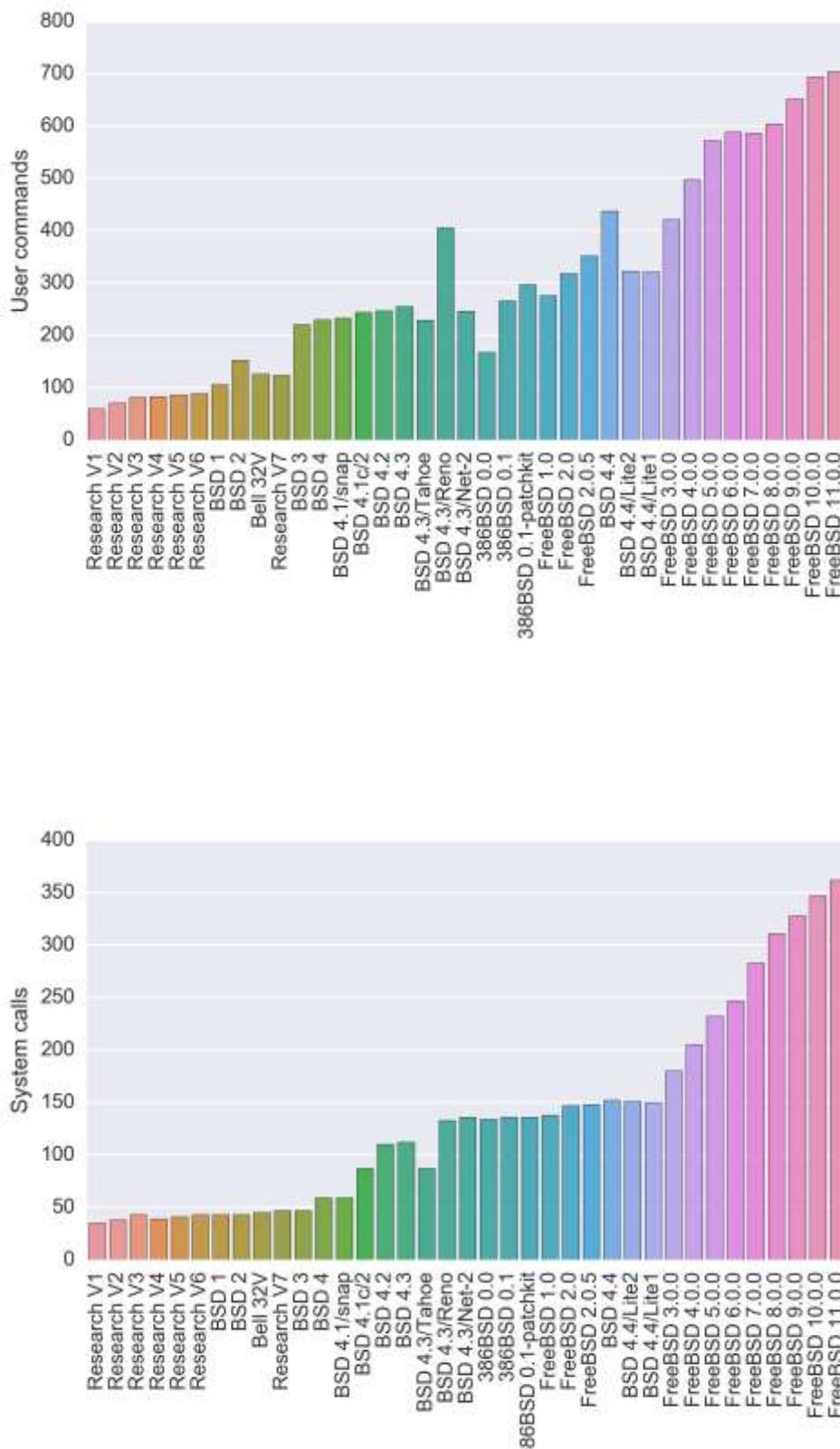


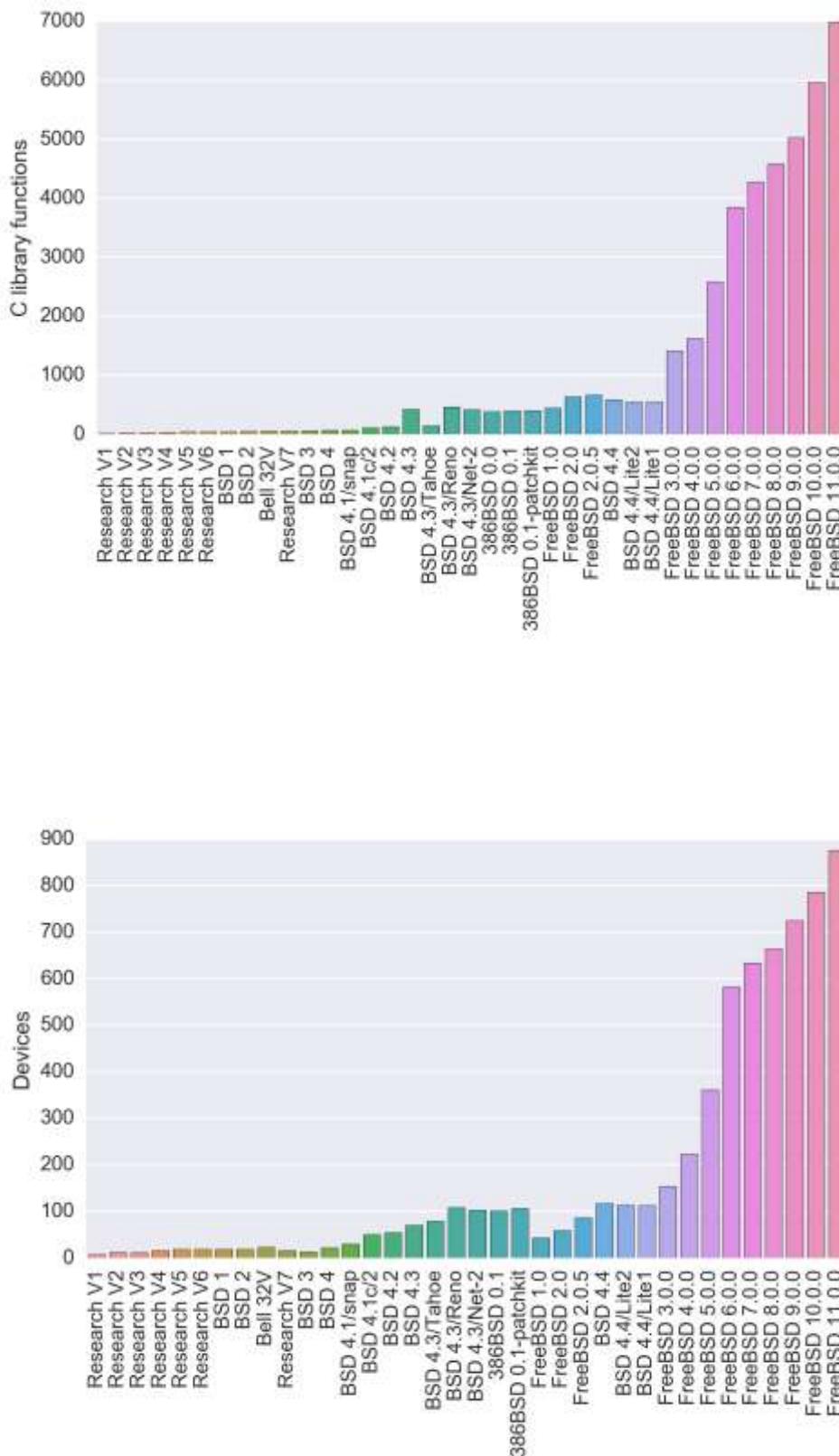
Agreement with Lehman's Laws

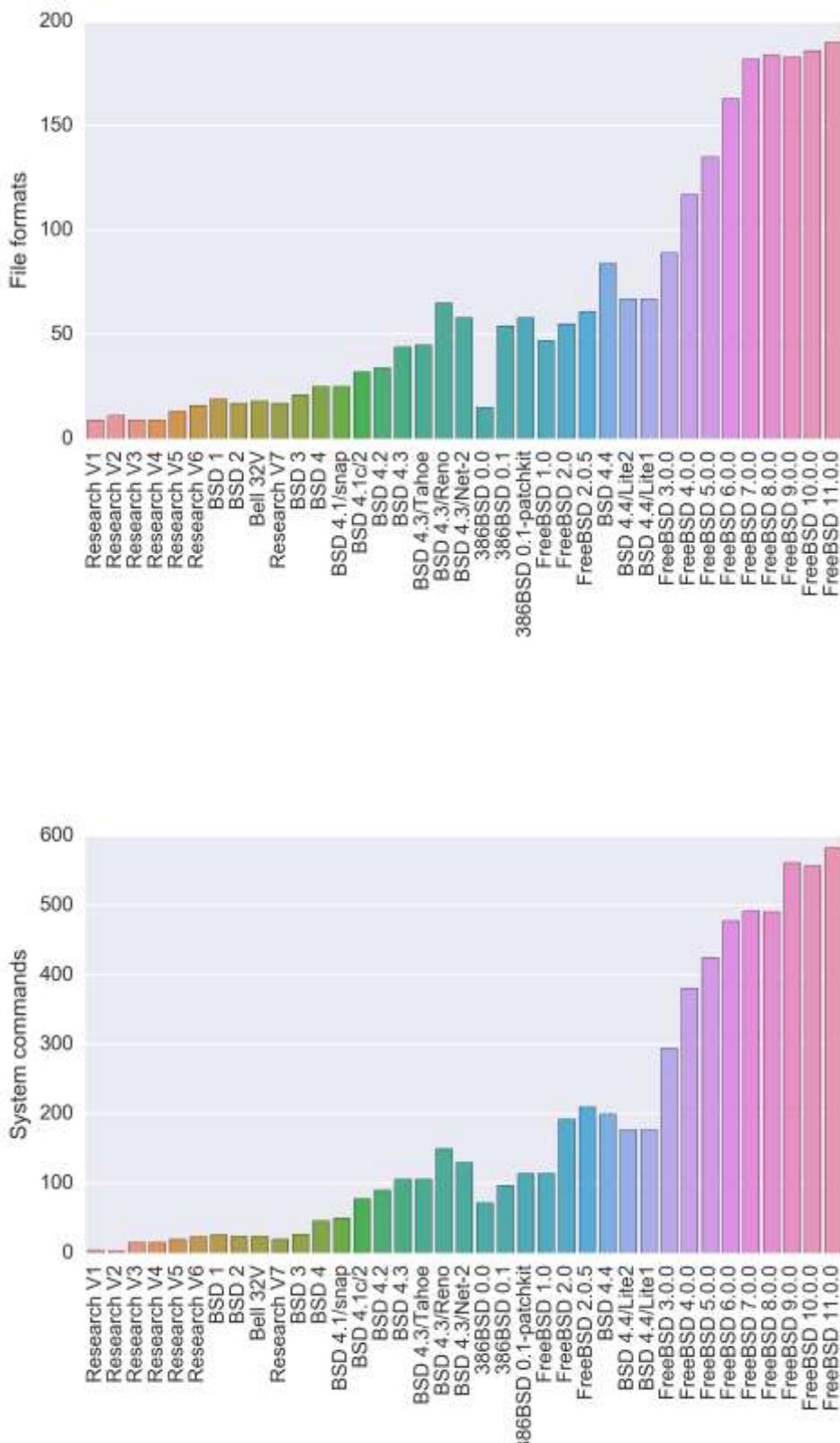
- Increasing Complexity
 - Conservation of Familiarity
 - Declining Quality
 - Feedback System

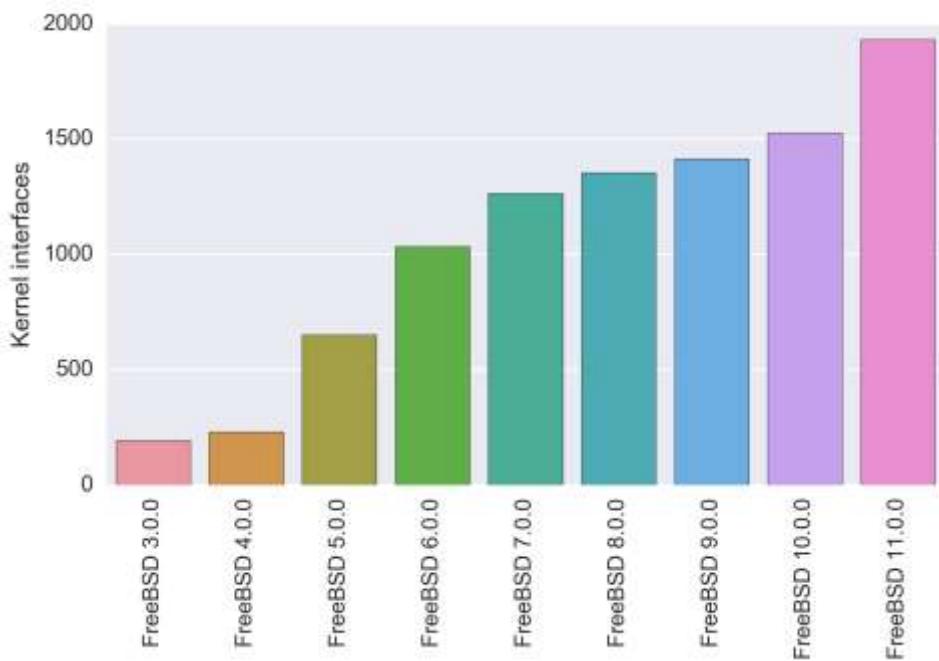


Growth in Facilities



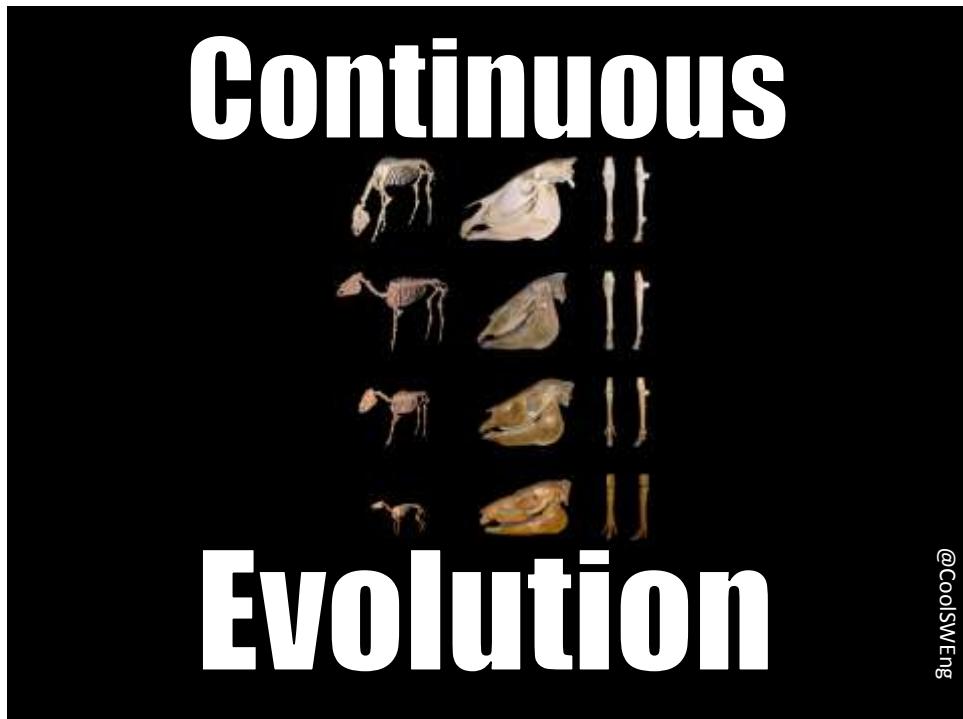


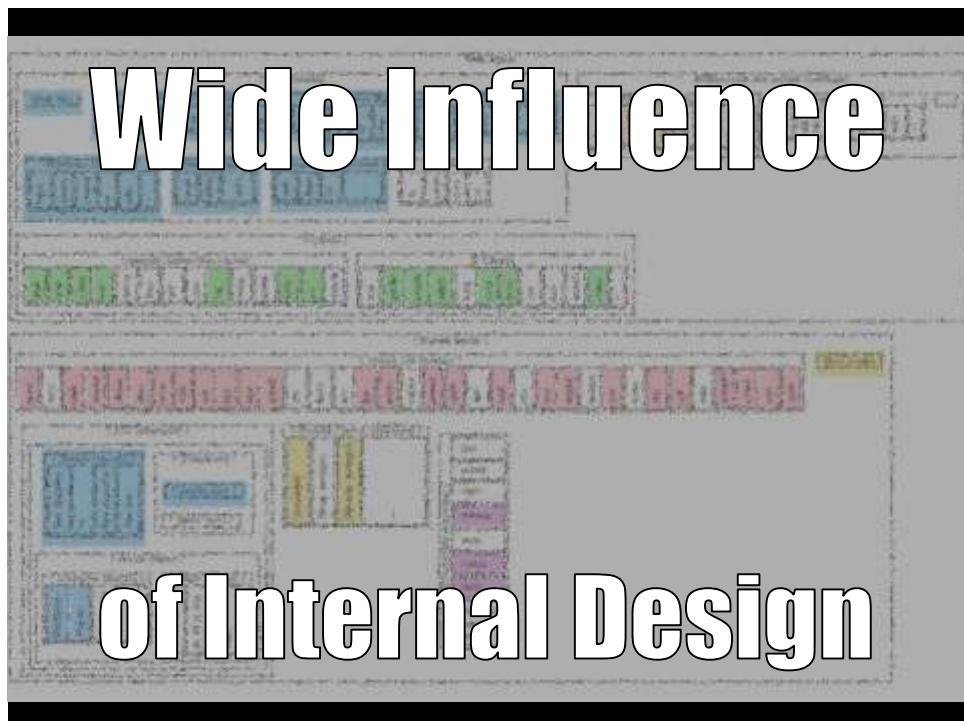
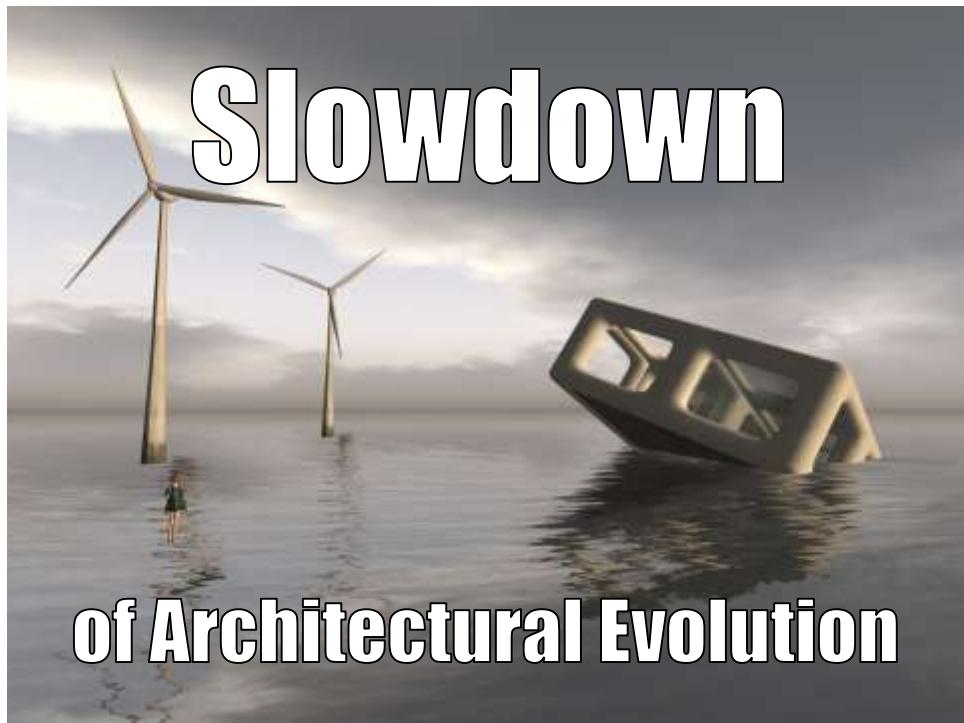


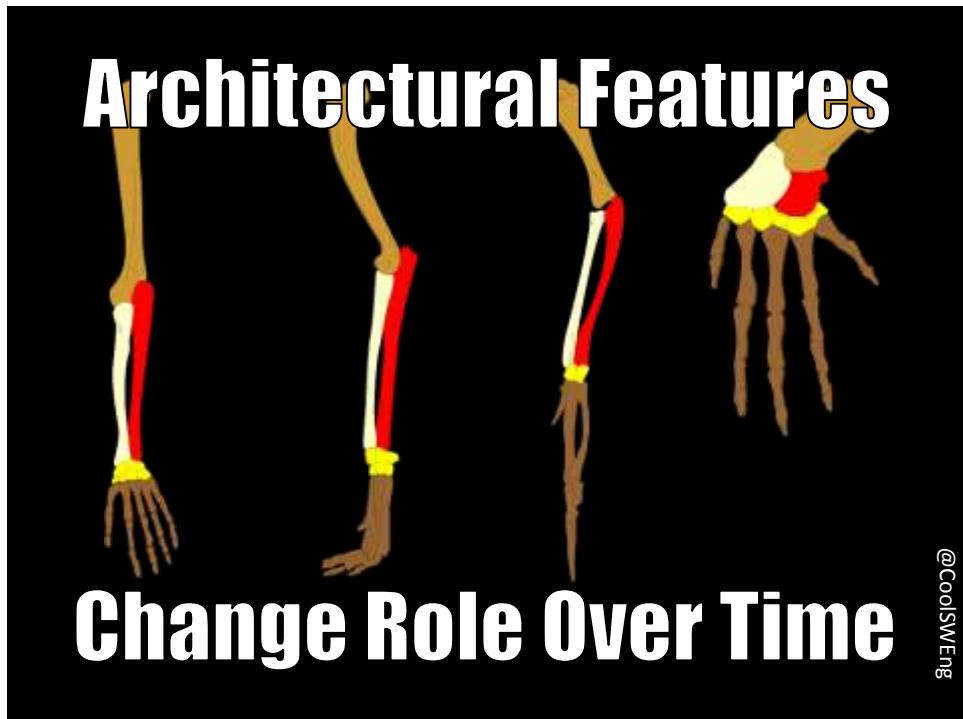




Conventions
Over Rigid Enforcement

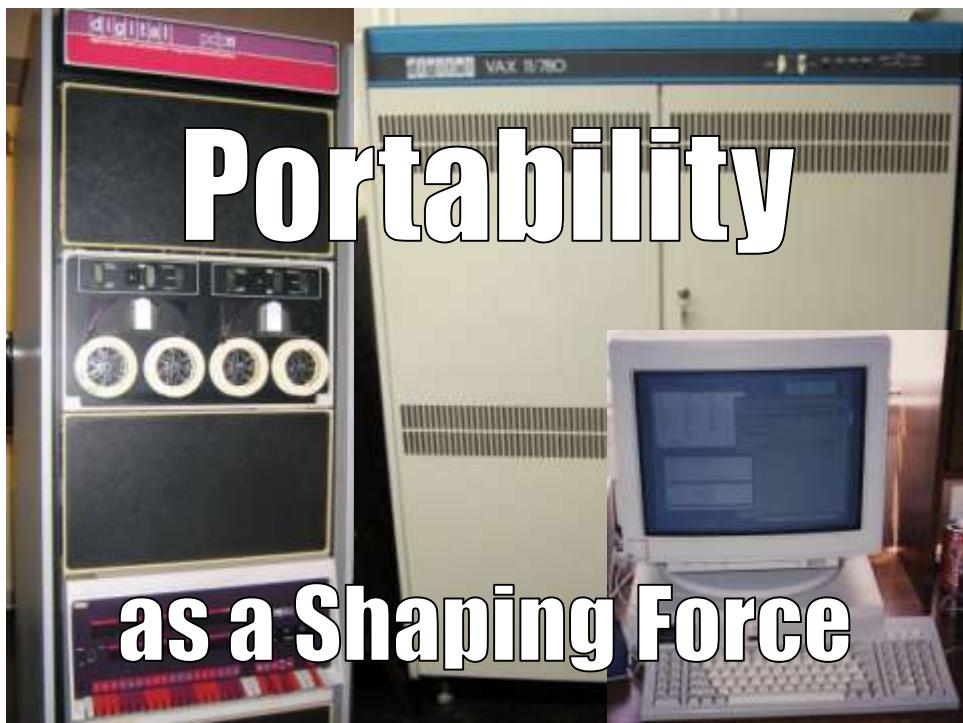


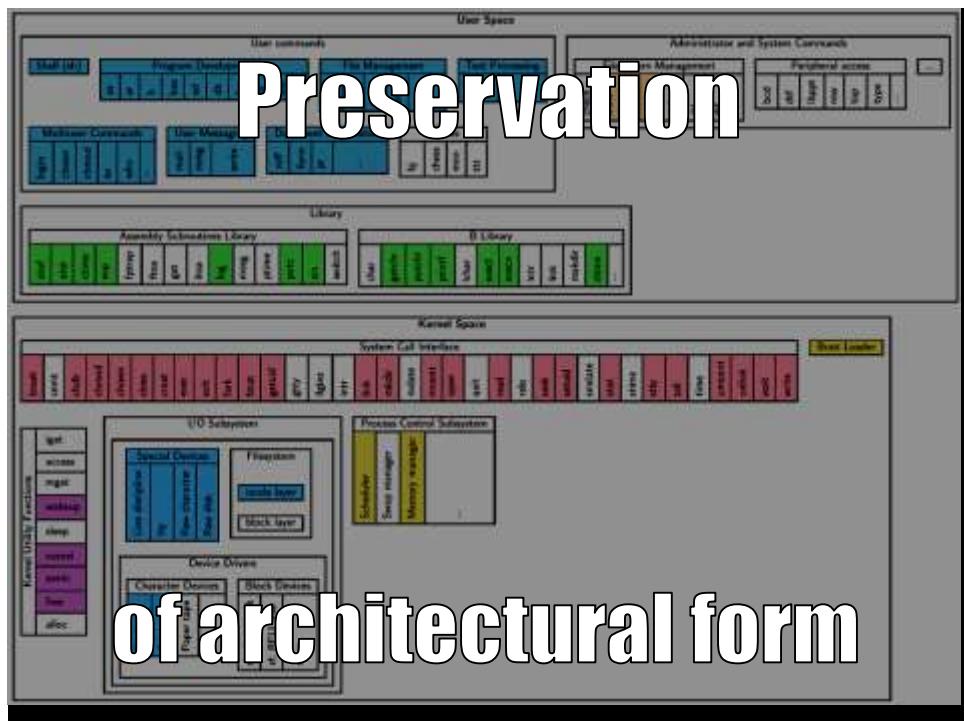


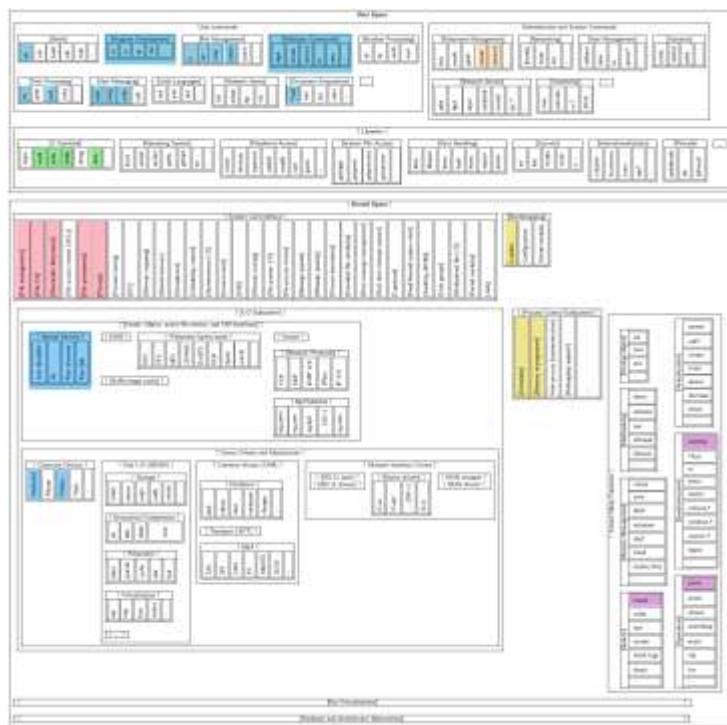
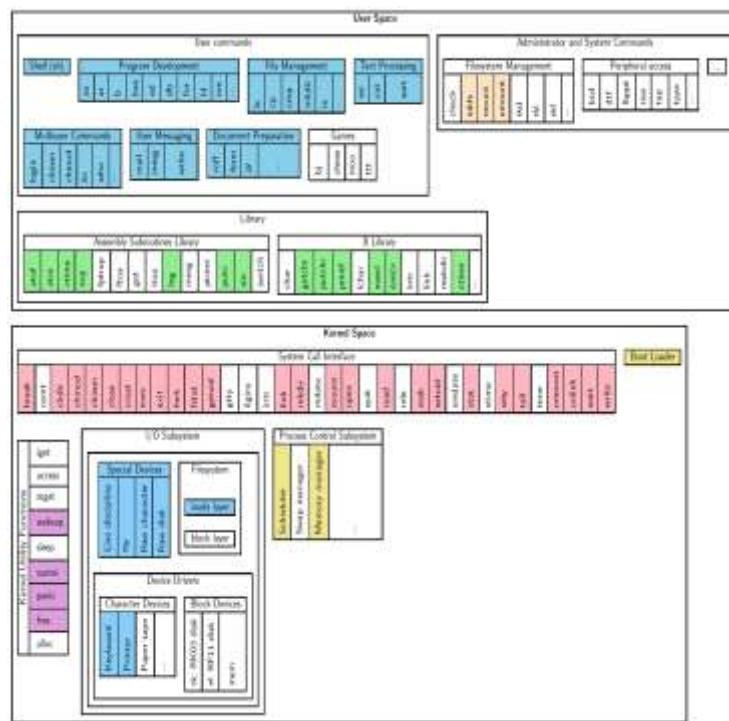


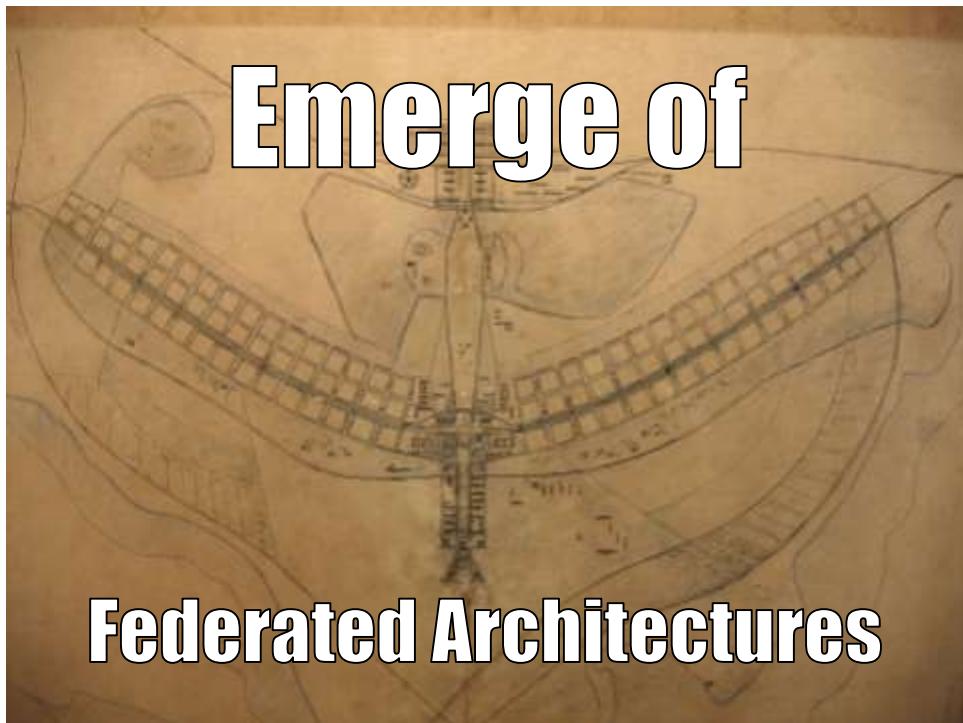
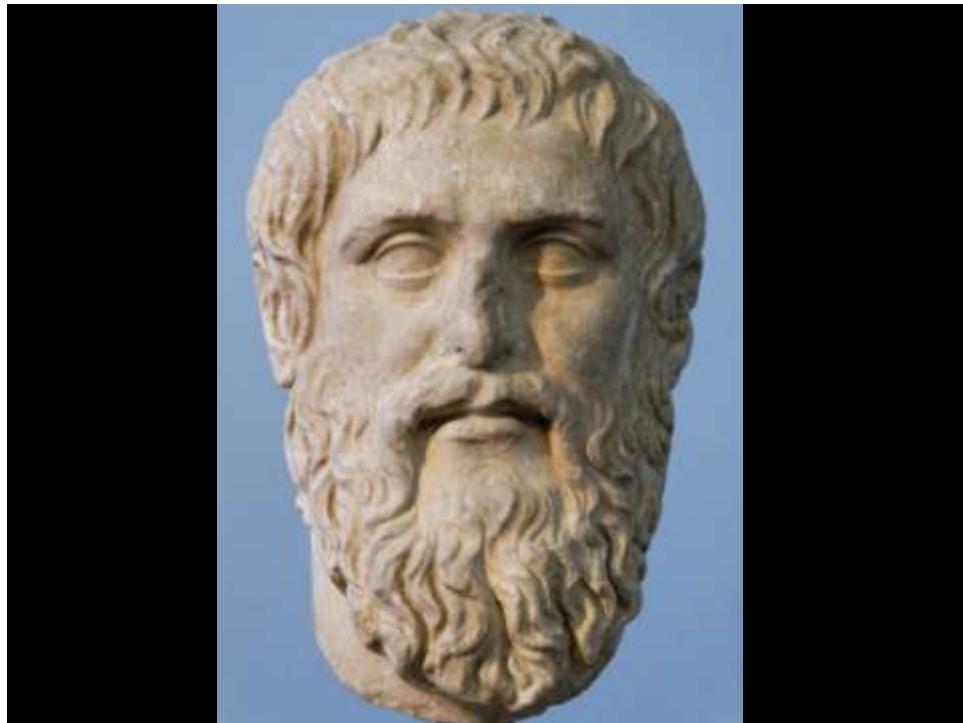


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Importance of a Gifted Team



Go and architect
something great!

Thank you!

github.com/dspinellis/unix-history-repo



www.spinellis.gr



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Funding Credit

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